

**D-OPTIMAL DESIGNS WITH
ORDERED CATEGORICAL DATA**

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Supplementary Materials

S.1 Commonly Used Link Functions for Cumulative Link Models

Link function	$g(\gamma)$	$g^{-1}(\eta)$	$(g^{-1})'(\eta)$
logit	$\log\left(\frac{\gamma}{1-\gamma}\right)$	$\frac{e^\eta}{1+e^\eta}$	$\frac{e^\eta}{(1+e^\eta)^2}$
probit	$\Phi^{-1}(\gamma)$	$\Phi(\eta)$	$\phi(\eta)$
log-log	$-\log[-\log(\gamma)]$	$\exp\{-e^{-\eta}\}$	$\exp\{-\eta - e^{-\eta}\}$
c-log-log	$\log[-\log(1 - \gamma)]$	$1 - \exp\{-e^\eta\}$	$\exp\{\eta - e^\eta\}$
cauchit	$\tan[\pi(\gamma - \frac{1}{2})]$	$\frac{1}{\pi} \arctan(\eta) + \frac{1}{2}$	$\frac{1}{\pi(1+\eta^2)}$

where $\Phi^{-1}(\cdot)$ is the cumulative distribution function of $N(0, 1)$, $\phi(\cdot)$ is the probability density function of $N(0, 1)$, and “c-log-log” stands for complementary log-log.

Example 1 (*continued*) For logit link g , $g^{-1}(\eta) = e^\eta/(1+e^\eta)$ and $(g^{-1})' = g^{-1}(1 - g^{-1})$. Thus $g_{ij} = (g^{-1})'(\theta_j - \mathbf{x}_i^T \boldsymbol{\beta}) = \gamma_{ij}(1 - \gamma_{ij})$. With $J = 3$, we have $\pi_{i1} + \pi_{i2} + \pi_{i3} = 1$ for $i = 1, \dots, m$. Then for $i = 1, \dots, m$, $g_{i1} = \pi_{i1}(\pi_{i2} + \pi_{i3})$, $g_{i2} = (\pi_{i1} + \pi_{i2})\pi_{i3}$, $b_{i2} = \pi_{i1}\pi_{i3}\pi_{i2}^{-1}(\pi_{i1} + \pi_{i2})(\pi_{i2} + \pi_{i3})$, $u_{i1} = \pi_{i1}\pi_{i2}^{-1}(\pi_{i1} + \pi_{i2})(\pi_{i2} + \pi_{i3})^2$, $u_{i2} = \pi_{i3}\pi_{i2}^{-1}(\pi_{i1} + \pi_{i2})^2(\pi_{i2} + \pi_{i3})$, $c_{i1} = \pi_{i1}(\pi_{i1} + \pi_{i2})(\pi_{i2} + \pi_{i3})$, $c_{i2} = \pi_{i3}(\pi_{i1} + \pi_{i2})(\pi_{i2} + \pi_{i3})$, $e_i = (\pi_{i1} + \pi_{i2})(\pi_{i1} + \pi_{i3})(\pi_{i2} + \pi_{i3})$. \square

S.2 Additional Lemmas

For Section 2: Since $(Y_{i1}, \dots, Y_{iJ}), i = 1, \dots, m$ are m independent random vectors, the log-likelihood function (up to a constant) of the cumulative link model is

$$l(\beta_1, \dots, \beta_d, \theta_1, \dots, \theta_{J-1}) = \sum_{i=1}^m \sum_{j=1}^J Y_{ij} \log(\pi_{ij})$$

where $\pi_{ij} = \gamma_{ij} - \gamma_{i,j-1}$ with $\gamma_{ij} = g^{-1}(\theta_j - \mathbf{x}_i^T \boldsymbol{\beta})$ for $j = 1, \dots, J-1$ and $\gamma_{i0} = 0, \gamma_{iJ} = 1, i = 1, \dots, m$. For $s = 1, \dots, d, t = 1, \dots, J-1$,

$$\begin{aligned} \frac{\partial l}{\partial \beta_s} &= \sum_{i=1}^m (-x_{is}) \cdot \left\{ \frac{Y_{i1}}{\pi_{i1}} \cdot (g^{-1})'(\theta_1 - \mathbf{x}_i^T \boldsymbol{\beta}) \right. \\ &\quad + \frac{Y_{i2}}{\pi_{i2}} \cdot [(g^{-1})'(\theta_2 - \mathbf{x}_i^T \boldsymbol{\beta}) - (g^{-1})'(\theta_1 - \mathbf{x}_i^T \boldsymbol{\beta})] \\ &\quad \left. + \dots + \frac{Y_{iJ}}{\pi_{iJ}} [-(g^{-1})'(\theta_{J-1} - \mathbf{x}_i^T \boldsymbol{\beta})] \right\} \\ \frac{\partial l}{\partial \theta_t} &= \sum_{i=1}^m (g^{-1})'(\theta_t - \mathbf{x}_i^T \boldsymbol{\beta}) \left(\frac{Y_{it}}{\pi_{it}} - \frac{Y_{i,t+1}}{\pi_{i,t+1}} \right) \end{aligned}$$

Since Y_{ij} 's come from multinomial distributions, we know $E(Y_{ij}) = n_i \pi_{ij}$, $E(Y_{ij}^2) = n_i(n_i - 1)\pi_{ij}^2 + n_i \pi_{ij}$, and $E(Y_{is}Y_{it}) = n_i(n_i - 1)\pi_{is}\pi_{it}$ when $s \neq t$. Then we have the following lemma:

Lemma S.1. *Let $\mathbf{F} = (F_{st})$ be the $(d+J-1) \times (d+J-1)$ Fisher information matrix.*

(i) For $1 \leq s \leq d, 1 \leq t \leq d$,

$$F_{st} = E \left(\frac{\partial l}{\partial \beta_s} \frac{\partial l}{\partial \beta_t} \right) = \sum_{i=1}^m n_i x_{is} x_{it} \sum_{j=1}^J \frac{(g_{ij} - g_{i,j-1})^2}{\pi_{ij}}$$

where $g_{ij} = (g^{-1})'(\theta_j - \mathbf{x}_i^T \boldsymbol{\beta}) > 0$ for $j = 1, \dots, J-1$ and $g_{i0} = g_{iJ} = 0$.

(ii) For $1 \leq s \leq d, 1 \leq t \leq J-1$,

$$F_{s,d+t} = E \left(\frac{\partial l}{\partial \beta_s} \frac{\partial l}{\partial \theta_t} \right) = \sum_{i=1}^m n_i (-x_{is}) g_{it} \left(\frac{g_{it} - g_{i,t-1}}{\pi_{it}} - \frac{g_{i,t+1} - g_{it}}{\pi_{i,t+1}} \right)$$

(iii) For $1 \leq s \leq J-1, 1 \leq t \leq d$,

$$F_{d+s,t} = E \left(\frac{\partial l}{\partial \theta_s} \frac{\partial l}{\partial \beta_t} \right) = \sum_{i=1}^m n_i (-x_{it}) g_{is} \left(\frac{g_{is} - g_{i,s-1}}{\pi_{is}} - \frac{g_{i,s+1} - g_{is}}{\pi_{i,s+1}} \right)$$

(iv) For $1 \leq s \leq J-1, 1 \leq t \leq J-1$,

$$F_{d+s,d+t} = E \left(\frac{\partial l}{\partial \theta_s} \frac{\partial l}{\partial \theta_t} \right) = \begin{cases} \sum_{i=1}^m n_i g_{is}^2 (\pi_{is}^{-1} + \pi_{i,s+1}^{-1}), & \text{if } s = t \\ \sum_{i=1}^m n_i g_{is} g_{it} (-\pi_{i,s \vee t}^{-1}), & \text{if } |s - t| = 1 \\ 0, & \text{if } |s - t| \geq 2 \end{cases}$$

where $s \vee t = \max\{s, t\}$.

Perevozskaya et al. (2003) obtained a detailed form of Fisher information matrix for logit link and one predictor. Our expressions here are good for fairly general link and d predictors. To simplify the notations, we denote for $i = 1, \dots, m$,

$$e_i = \sum_{j=1}^J \frac{(g_{ij} - g_{i,j-1})^2}{\pi_{ij}} > 0 \quad (\text{S.1})$$

$$c_{it} = g_{it} \left(\frac{g_{it} - g_{i,t-1}}{\pi_{it}} - \frac{g_{i,t+1} - g_{it}}{\pi_{i,t+1}} \right), \quad t = 1, \dots, J-1 \quad (\text{S.2})$$

$$u_{it} = g_{it}^2 (\pi_{it}^{-1} + \pi_{i,t+1}^{-1}) > 0, \quad t = 1, \dots, J-1 \quad (\text{S.3})$$

$$b_{it} = g_{i,t-1} g_{it} \pi_{it}^{-1} > 0, \quad t = 2, \dots, J-1 \text{ (if } J \geq 3) \quad (\text{S.4})$$

Note that g_{ij} is defined in Lemma S.1 (i). Then we obtain the following lemma which plays a key role in calculating $|\mathbf{F}|$.

Lemma S.2. $c_{it} = u_{it} - b_{it} - b_{i,t+1}$, $i = 1, \dots, m$; $t = 1, \dots, J-1$; $e_i = \sum_{t=1}^{J-1} c_{it} = \sum_{t=1}^{J-1} (u_{it} - 2b_{it})$, $i = 1, \dots, m$, where $b_{i1} = b_{iJ} = 0$ for $i = 1, \dots, m$.

Lemma S.3. $\text{Rank}((\mathbf{A}_{i1} \mathbf{A}_{i2})) \leq 1$ where “=” is true if and only if $\mathbf{x}_i \neq 0$.

Based on Lemmas 1 and S.3, we obtain the two lemmas below on $c_{\alpha_1, \dots, \alpha_m}$ which significantly simplify the structure of $|\mathbf{F}|$ as a polynomial of (n_1, \dots, n_m) .

Lemma S.4. If $\max_{1 \leq i \leq m} \alpha_i \geq J$, then $|\mathbf{A}_\tau| = 0$ for any $\tau \in (\alpha_1, \dots, \alpha_m)$ and thus $c_{\alpha_1, \dots, \alpha_m} = 0$.

Proof of Lemma S.4: Without any loss of generality, we assume $\alpha_1 \geq \alpha_2 \geq \dots \geq \alpha_m$. Then $\max_{1 \leq i \leq m} \alpha_i \geq J$ implies $\alpha_1 \geq J$. In this case, for any $\tau \in (\alpha_1, \dots, \alpha_m)$, $\tau^{-1}(1) := \{i \mid \tau(i) = 1\} \subset \{1, \dots, d+J-1\}$ and $|\tau^{-1}(1)| = \alpha_1$. If $|\tau^{-1}(1) \cap \{1, \dots, d\}| \geq 2$, then $|\mathbf{A}_\tau| = 0$ due to Lemma S.3; otherwise $\{d+1, \dots, d+J-1\} \subset \tau^{-1}(1)$ and thus $|\mathbf{A}_\tau| = 0$ due to Lemma 1. Thus $c_{\alpha_1, \dots, \alpha_m} = 0$ according to (2.3) provided in Theorem 2. \square

Lemma S.5. If $\#\{i : \alpha_i \geq 1\} \leq d$, then $|\mathbf{A}_\tau| = 0$ for any $\tau \in (\alpha_1, \dots, \alpha_m)$ and thus $c_{\alpha_1, \dots, \alpha_m} = 0$.

Proof of Lemma S.5: Without any loss of generality, we assume $\alpha_1 \geq \alpha_2 \geq \dots \geq \alpha_m$. Then $\#\{i : \alpha_i \geq 1\} \leq d$ indicates $\alpha_{d+1} = \dots = \alpha_m = 0$. Let $\tau : \{1, 2, \dots, d+J-1\} \rightarrow \{1, \dots, m\}$ satisfy $\tau \in (\alpha_1, \dots, \alpha_m)$. Then the $(d+J-1) \times (d+J-1)$ matrix \mathbf{A}_τ can be written as

$$\begin{pmatrix} A_{\tau 1} & A_{\tau 2} \\ A_{\tau 3} & A_{\tau 4} \end{pmatrix}$$

$$= \begin{pmatrix} (e_{\tau(s)}x_{\tau(s)}s x_{\tau(s)}t)_{s=1,\dots,d;t=1,\dots,d} & (-x_{\tau(s)}s c_{\tau(s)}t)_{s=1,\dots,d;t=1,\dots,J-1} \\ (-c_{\tau(d+s)}s x_{\tau(d+s)}t)_{s=1,\dots,J-1;t=1,\dots,d} & A_{\tau_4} \end{pmatrix}$$

where the $(J-1) \times (J-1)$ matrix A_{τ_4} is either a single entry $u_{\tau(d+1)1}$ (if $J = 2$) or symmetric tri-diagonal with diagonal entries $u_{\tau(d+1)1}, \dots, u_{\tau(d+J-1),J-1}$, upper off-diagonal entries $-b_{\tau(d+1)2}, \dots, -b_{\tau(d+J-2),J-1}$, and lower off-diagonal entries $-b_{\tau(d+2)2}, \dots, -b_{\tau(d+J-1),J-1}$. Note that A_{τ} is asymmetric in general.

If $\#\{i : \alpha_i \geq 1\} \leq d-1$, then there exists an i_0 such that $1 \leq i_0 \leq d$ and $|\tau^{-1}(i_0) \cap \{1, \dots, d\}| \geq 2$. In this case, $|A_{\tau}| = 0$ according to Lemma S.3.

If $\#\{i : \alpha_i \geq 1\} = d$, we may assume $|\tau^{-1}(i) \cap \{1, \dots, d\}| = 1$ for $i = 1, \dots, d$ (otherwise $|A_{\tau}| = 0$ according to Lemma S.3). Suppose $\alpha_1 \geq \alpha_2 \geq \dots \geq \alpha_k \geq 2 > \alpha_{k+1}$. Then $\{d+1, \dots, d+J-1\} \subset \cup_{i=1}^k \tau^{-1}(i)$ and $\sum_{i=1}^k (\alpha_i - 1) = J-1$. In order to show $|A_{\tau}| = 0$, we first replace A_{τ_1} with $A_{\tau_1}^{(1)} = (e_{\tau(s)}x_{\tau(s)}t)_{s=1,\dots,d;t=1,\dots,d}$ and replace A_{τ_2} with $A_{\tau_2}^{(1)} = (-c_{\tau(s)}t)_{s=1,\dots,d;t=1,\dots,J-1}$. It changes A_{τ} into a new matrix $A_{\tau}^{(1)}$. Note that $|A_{\tau}| = \prod_{s=1}^d x_{\tau(s)}s \cdot |A_{\tau}^{(1)}|$. According to Lemma S.2, the sum of the columns of $A_{\tau_2}^{(1)}$ is $(-e_{\tau(1)}, \dots, -e_{\tau(d)})^T$, and the elementwise sum of the columns of A_{τ_4} is $(c_{\tau(d+1)1}, c_{\tau(d+2)2}, \dots, c_{\tau(d+J-1),J-1})^T$. Secondly, for $t = 1, \dots, d$, we add $x_{1t}(-e_{\tau(1)}, \dots, -e_{\tau(d)}, c_{\tau(d+1)1}, \dots, c_{\tau(d+J-1),J-1})^T$ to the t th column of $A_{\tau}^{(1)}$. We denote the resulting matrix by $A_{\tau}^{(2)}$. Note that $|A_{\tau}^{(1)}| = |A_{\tau}^{(2)}|$. We consider the sub-matrix $A_{\tau_d}^{(2)}$ which consists of the first d columns of $A_{\tau}^{(2)}$. For $s \in \tau^{-1}(1)$, the s th row of $A_{\tau_d}^{(2)}$ is simply 0. For $i = 2, \dots, k$, the j th row of $A_{\tau_d}^{(2)}$ is proportional to $(x_{i1} - x_{11}, x_{i2} - x_{12}, \dots, x_{id} - x_{1d})$ if $j \in \tau^{-1}(i)$. Therefore, $\text{Rank}(A_{\tau_d}^{(2)}) \leq (d+J-1) - \alpha_1 - \sum_{i=2}^k (\alpha_i - 1) = d-1$, which leads to $|A_{\tau_d}^{(2)}| = 0$ and thus $|A_{\tau}^{(1)}| = 0$, $|A_{\tau}| = 0$. According to (2.3) in Theorem 2, $c_{\alpha_1, \dots, \alpha_m} = 0$. \square

Lemma S.6. $\mathbf{F} = \mathbf{F}(\mathbf{p})$ is always positive semi-definite. It is positive definite if and only if $\mathbf{p} \in S_+$. Furthermore, $\log f(\mathbf{p})$ is concave on S .

For Section 5.2: The procedure seeking for analytic solutions here follows Tong, Volkmer, and Yang (2014). As a direct conclusion of the Karush-Kuhn-Tucker conditions (see also Theorem 10), a necessary condition for (p_1, p_2, p_3) to maximize $f(p_1, p_2, p_3)$ in (5.5) is (5.6), which are equivalent to $\partial f / \partial p_1 = \partial f / \partial p_3$ and $\partial f / \partial p_2 = \partial f / \partial p_3$. In terms of p_i, w_i 's, they are

$$(p_3 - p_1)(p_1 w_1 + p_2 w_2 + p_3 w_3) = (w_3 - w_1)p_1 p_3 \quad (\text{S.5})$$

$$(p_3 - p_2)(p_1 w_1 + p_2 w_2 + p_3 w_3) = (w_3 - w_2)p_2 p_3 \quad (\text{S.6})$$

Denote $y_1 = p_1/p_3 > 0$ and $y_2 = p_2/p_3 > 0$. Since $p_1 + p_2 + p_3 = 1$, it implies $p_3 = 1/(y_1 + y_2 + 1)$, $p_1 = y_1/(y_1 + y_2 + 1)$, and $p_2 = y_2/(y_1 + y_2 + 1)$. In terms of y_1, y_2 , (S.5) and (S.6) are equivalent to

$$(1 - y_1)(y_1 w_1 + y_2 w_2 + w_3) = (w_3 - w_1)y_1 \quad (\text{S.7})$$

$$(1 - y_2)(y_1 w_1 + y_2 w_2 + w_3) = (w_3 - w_2)y_2 \quad (\text{S.8})$$

Lemma S.7. *Suppose $0 < w_3 < w_2 < w_1$. If (p_1, p_2, p_3) maximizes $f(p_1, p_2, p_3)$ in (5.5) under the constraints $p_1, p_2, p_3 \geq 0$ and $p_1 + p_2 + p_3 = 1$, then $0 < p_3 \leq p_2 \leq p_1 < 1$.*

The proof of the lemma above is straightforward, because otherwise one could exchange p_i, p_j to strictly improve $f(p_1, p_2, p_3)$. Now we are ready to get solutions to equations (S.7) and (S.8) case by case.

- (i) $w_1 = w_3$. In that case, (S.7) implies $y_1 = 1$. After plugging it into (S.8), the only positive solution is $y_2 = (-3w_1 + 2w_2 + \sqrt{9w_1^2 - 4w_1w_2 + 4w_2^2})/(2w_2)$.
- (ii) $w_2 = w_3$. In that case, (S.8) implies $y_2 = 1$. After plugging it into (S.7), the only positive solution is $y_1 = (2w_1 - 3w_2 + \sqrt{4w_1^2 - 4w_1w_2 + 9w_2^2})/(2w_1)$.
- (iii) $w_1 = w_2$ but $w_1 \neq w_3$. The ratio of (S.7) and (S.8) leads to $y_1 = y_2$. After plugging it into (S.7), the only positive solution is $y_1 = (3w_1 - 2w_3 + \sqrt{9w_1^2 - 4w_1w_3 + 4w_3^2})/(4w_1)$.
- (iv) w_1, w_2, w_3 are distinct. Without any loss of generality, we assume $0 < w_3 < w_2 < w_1$, because otherwise the previous elimination procedure in the order of p_3, p_2, p_1 could be easily changed accordingly. Based on Lemma S.7, if (p_1, p_2, p_3) maximizes f_4 , then $0 < p_3 \leq p_2 \leq p_1 < 1$ and thus $y_1 \geq y_2 \geq 1$. The ratio of (S.7) and (S.8) leads to $(1 - y_1)/(1 - y_2) = (w_3 - w_1)/(w_3 - w_2) \cdot y_1/y_2$, which implies

$$y_2 = \frac{(w_1 - w_3)y_1}{(w_2 - w_3) + (w_1 - w_2)y_1}. \quad (\text{S.9})$$

Note that $(w_2 - w_3) + (w_1 - w_2)y_1 \geq w_1 - w_3 > 0$. After plugging (S.9) into (S.7), we get

$$c_0 + c_1 y_1 + c_2 y_1^2 + c_3 y_1^3 = 0 \quad (\text{S.10})$$

where $c_0 = w_3(w_2 - w_3) > 0$, $c_1 = 3w_1w_2 - w_1w_3 - 4w_2w_3 + 2w_3^2 > 0$, $c_2 = 2w_1^2 - 4w_1w_2 - w_1w_3 + 3w_2w_3$, $c_3 = w_1(w_2 - w_1) < 0$.

Lemma S.8. *Suppose $0 < w_3 < w_2 < w_1$. Then equation (S.10) has one and only one solution $y_1^* \geq 1$. Furthermore, $y_1^* > 1$.*

Proof of Lemma S.8: In order to locate the roots of equation (S.10), we let $f_1(y_1) = c_0 + c_1y_1 + c_2y_1^2 + c_3y_1^3$. Then $f_1(1) = c_0 + c_1 + c_2 + c_3 = (w_1 - w_3)^2 > 0$.

On the other hand, the first derivative of f_1 is $f_1'(y_1) = a_0 + a_1y_1 + a_2y_1^2$, where $a_0 = 3w_1w_2 - w_1w_3 - 4w_2w_3 + 2w_3^2 = w_1(w_2 - w_3) + 2(w_1 - w_2)w_2 + 2(w_2 - w_3)^2 > 0$, $a_1 = 2(2w_1^2 - 4w_1w_2 - w_1w_3 + 3w_2w_3)$, and $a_2 = 3w_1(w_2 - w_1) < 0$. Therefore, $a_1^2 - 4a_0a_2 > a_1^2 \geq 0$ and $f_1'(y_1) = a_2(y_1 - y_{11})(y_1 - y_{12})$, where

$$y_{11} = \frac{-a_1 + \sqrt{a_1^2 - 4a_0a_2}}{2a_2} < 0, \quad y_{12} = \frac{-a_1 - \sqrt{a_1^2 - 4a_0a_2}}{2a_2} > y_{11}$$

It can be verified that $y_{12} < 1$ if and only if $w_1 < 2(w_2 + w_3)$. There are two cases: *Case (i):* If $y_{12} < 1$, then $f_1'(y_1) < 0$ for all $y_1 > 1$. That is, $f_1(y_1)$ strictly decreases after $y_1 = 1$. Since $f_1(1) > 0$ and $f_1(\infty) = -\infty$, then there is one and only one solution in $(1, \infty)$; *Case (ii):* If $y_{12} \geq 1$, then $f_1'(y_1) \geq 0$ for $y_1 \in [1, y_{12}]$ and $f_1'(y_1) < 0$ for $y_1 \in (y_{12}, \infty)$. That is, $f_1(y_1)$ increases in $[1, y_{12}]$ and then strictly decreases in (y_{12}, ∞) . Again, due to $f_1(1) > 0$ and $f_1(\infty) = -\infty$, there is one and only one solution in $(1, \infty)$. In either case, the conclusion is justified. \square

S.3 Additional Proofs

Proof of Theorem 1 It is a direct conclusion of Lemmas S.1 and S.2. \square

Examples of \mathbf{A}_{i3} in Theorem 1 include (u_{i1}) ,

$$\begin{pmatrix} u_{i1} & -b_{i2} \\ -b_{i2} & u_{i2} \end{pmatrix}, \begin{pmatrix} u_{i1} & -b_{i2} & 0 \\ -b_{i2} & u_{i2} & -b_{i3} \\ 0 & -b_{i3} & u_{i3} \end{pmatrix}, \begin{pmatrix} u_{i1} & -b_{i2} & 0 & 0 \\ -b_{i2} & u_{i2} & -b_{i3} & 0 \\ 0 & -b_{i3} & u_{i3} & -b_{i4} \\ 0 & 0 & -b_{i4} & u_{i4} \end{pmatrix}$$

for $J = 2, 3, 4$, or 5 respectively.

Proof of Theorem 2 To study the structure of $|\mathbf{F}|$ as a polynomial function of (n_1, \dots, n_m) , we denote the (k, l) th entry of \mathbf{A}_i by $a_{kl}^{(i)}$. Given a row map $\tau : \{1, 2, \dots, d + J - 1\} \rightarrow \{1, \dots, m\}$, we define a $(d + J - 1) \times (d + J - 1)$ matrix $\mathbf{A}_\tau = \left(a_{kl}^{(\tau(k))} \right)$ whose k th row is given by the k th row of

$\mathbf{A}_{\tau^{(k)}}$. For a power index $(\alpha_1, \dots, \alpha_m)$ with $\alpha_i \in \{0, 1, \dots, d + J - 1\}$ and $\sum_{i=1}^m \alpha_i = d + J - 1$, we denote

$$\tau \in (\alpha_1, \dots, \alpha_m)$$

if $\alpha_i = \#\{j : \tau(j) = i\}$ for each $i = 1, \dots, m$. In terms of the construction of \mathbf{A}_{τ} , it says that α_i rows of \mathbf{A}_{τ} are from the matrix \mathbf{A}_i .

According to the Leibniz formula for the determinant,

$$|\mathbf{F}| = \left| \sum_{i=1}^m n_i \mathbf{A}_i \right| = \sum_{\sigma \in S_{d+J-1}} (-1)^{\text{sgn}(\sigma)} \prod_{k=1}^{d+J-1} \sum_{i=1}^m n_i a_{k, \sigma(k)}^{(i)}$$

where σ is a permutation of $\{1, 2, \dots, d + J - 1\}$, and $\text{sgn}(\sigma)$ is the sign or signature of σ . Therefore,

$$\begin{aligned} c_{\alpha_1, \dots, \alpha_m} &= \sum_{\sigma \in S_{d+J-1}} (-1)^{\text{sgn}(\sigma)} \sum_{\tau \in (\alpha_1, \dots, \alpha_m)} \prod_{k=1}^{d+J-1} a_{k, \sigma(k)}^{(\tau(k))} \\ &= \sum_{\tau \in (\alpha_1, \dots, \alpha_m)} \sum_{\sigma \in S_{d+J-1}} (-1)^{\text{sgn}(\sigma)} \prod_{k=1}^{d+J-1} a_{k, \sigma(k)}^{(\tau(k))} \\ &= \sum_{\tau \in (\alpha_1, \dots, \alpha_m)} |\mathbf{A}_{\tau}| \end{aligned}$$

□

Proof of Lemma 2 To simplify the notations, we let $i_s = s + 1$, $s = 0, \dots, d$. That is, $\alpha_1 = J - 1$, $\alpha_2 = \dots = \alpha_{d+1} = 1$. There are only two types of $\tau \in (\alpha_1, \dots, \alpha_m)$, such that, $|\mathbf{A}_{\tau}|$ may not be 0.

τ of type I: There exist $1 \leq k \leq d$, $2 \leq l \leq d + 1$, and $1 \leq q \leq J - 1$, such that, $\tau(k) = 1$ and $\tau(d + q) = l$. Following a similar procedure as in the proof of Lemma S.5, we obtain

$$|\mathbf{A}_{\tau}| = \prod_{i=2}^{d+1} e_i \cdot |\mathbf{A}_{13}| \cdot (-1)^d |\mathbf{X}_1[1, 2, \dots, d + 1]| \cdot (-1)^{\text{sgn}(\tau)} \prod_{s=1}^d x_{\tau(s)s} \cdot \frac{c_{lq}}{e_l}$$

τ of type II: $\tau(d + 1) = \dots = \tau(d + J - 1) = 1$ and $\{\tau(1), \dots, \tau(d)\} = \{2, \dots, d + 1\}$. It can be verified that

$$|\mathbf{A}_{\tau}| = \prod_{i=2}^{d+1} e_i \cdot |\mathbf{A}_{13}| \cdot (-1)^d |\mathbf{X}_1[1, 2, \dots, d + 1]| \cdot (-1)^{\text{sgn}(\tau)} \prod_{s=1}^d x_{\tau(s)s}$$

According to Theorem 2,

$$\begin{aligned}
c_{\alpha_1, \dots, \alpha_m} &= \sum_{\tau \text{ of type I}} |\mathbf{A}_\tau| + \sum_{\tau \text{ of type II}} |\mathbf{A}_\tau| \\
&= \prod_{i=2}^{d+1} e_i \cdot |\mathbf{A}_{13}| \cdot (-1)^d |\mathbf{X}_1[1, 2, \dots, d+1]| \cdot \left(\sum_{k=1}^d \sum_{l=2}^{d+1} \sum_{\tau \in S_{d+1}: \tau(k)=1, \tau(d+1)=l} \right. \\
&\quad \left. (-1)^{\text{sgn}(\tau)} \prod_{s=1}^d x_{\tau(s)s} \sum_{q=1}^{J-1} \frac{C_{lq}}{e_l} + \sum_{\tau \in S_{d+1}: \tau(d+1)=1} (-1)^{\text{sgn}(\tau)} \prod_{s=1}^d x_{\tau(s)s} \right) \\
&= \prod_{i=2}^{d+1} e_i \cdot |\mathbf{A}_{13}| \cdot (-1)^d |\mathbf{X}_1[1, 2, \dots, d+1]| \cdot \sum_{\tau \in S_{d+1}} (-1)^{\text{sgn}(\tau)} \prod_{s=1}^d x_{\tau(s)s} \\
&= \prod_{i=2}^{d+1} e_i \cdot |\mathbf{A}_{13}| \cdot (-1)^d |\mathbf{X}_1[1, 2, \dots, d+1]| \cdot (-1)^d |\mathbf{X}_1[1, 2, \dots, d+1]| \\
&= \prod_{i=2}^{d+1} e_i \cdot |\mathbf{A}_{13}| \cdot |\mathbf{X}_1[1, 2, \dots, d+1]|^2
\end{aligned}$$

where S_{d+1} is the set of permutations of $\{1, \dots, d+1\}$. The general case with i_0, i_1, \dots, i_d can be obtained similarly. \square

Proof of Theorem 4 Suppose $\text{Rank}(\mathbf{X}_1) = d+1$. Then there exist $i_0, \dots, i_d \in \{1, \dots, m\}$, such that, $|\mathbf{X}_1[i_0, i_1, \dots, i_d]| \neq 0$. According to Lemma S.4, $f(\mathbf{p})$ can be regarded as an order- $(J-1)$ polynomial of p_{i_0} . Let $p_{i_0} = x \in (0, 1)$ and $p_i = (1-x)/(m-1)$ for $i \neq i_0$. Based on Lemma 2, $f(\mathbf{p})$ can be written as

$$\begin{aligned}
f_{i_0}(x) &= a_{J-1} x^{J-1} \left(\frac{1-x}{m-1} \right)^d + a_{J-2} x^{J-2} \left(\frac{1-x}{m-1} \right)^{d+1} \\
&\quad + \dots + a_1 x \left(\frac{1-x}{m-1} \right)^{d+J-2} + a_0 \left(\frac{1-x}{m-1} \right)^{d+J-1}, \text{ where} \\
a_{J-1} &= |\mathbf{A}_{i_0 3}| \sum_{\{i'_1, \dots, i'_d\} \subset \{1, \dots, m\} \setminus \{i_0\}} \prod_{s=1}^d e_{i'_s} |\mathbf{X}_1[i_0, i'_1, \dots, i'_d]|^2 > 0
\end{aligned}$$

Therefore, $\lim_{x \rightarrow 1^-} (1-x)^{-d} x^{1-J} f_{i_0}(x) = (m-1)^{-d} a_{J-1} > 0$. That is, $f(\mathbf{p}) > 0$ for $p_{i_0} = x$ close enough to 1 and $p_i = (1-x)/(m-1)$ for $i \neq i_0$.

In order to justify that the condition $\text{Rank}(\mathbf{X}_1) = d+1$ is also necessary, we only need to show that $f(\mathbf{p}) \equiv 0$ if $\text{Rank}(\mathbf{X}_1) \leq d$. Actually, for any $\tau : \{1, \dots, d+J-1\} \rightarrow \{1, \dots, m\}$, we construct $\mathbf{A}_\tau^{(1)}$ as in the proof of Lemma S.5. Then $|\mathbf{A}_\tau| = \prod_{s=1}^d x_{\tau(s)s} \cdot |\mathbf{A}_\tau^{(1)}|$. Similar as in the proof of Lemma S.5, for $t = 1, \dots, d$, we add $x_{\tau(1)t}(-e_{\tau(1)}, \dots, -e_{\tau(d)}, c_{\tau(d+1)1}, \dots, c_{\tau(d+J-1), J-1})^T$ to the t th column of $\mathbf{A}_\tau^{(1)}$. We denote the resulting matrix by $\mathbf{A}_\tau^{(3)}$. Note that $|\mathbf{A}_\tau^{(1)}| = |\mathbf{A}_\tau^{(3)}|$. We consider the sub-matrix $\mathbf{A}_{\tau d}^{(3)}$ which consists of the first d columns of $\mathbf{A}_\tau^{(3)}$. For $s \in \tau^{-1}(\tau(1))$, the s th row of $\mathbf{A}_{\tau d}^{(3)}$ is simply 0. For $s = 2, \dots, k$, the s th row of $\mathbf{A}_{\tau d}^{(3)}$ is $e_{\tau(s)}(x_{\tau(s)1} - x_{\tau(1)1}, \dots, x_{\tau(s)d} - x_{\tau(1)d})$. For $s = 1, \dots, J-1$, the $(d+s)$ th row of $\mathbf{A}_{\tau d}^{(3)}$ is $-c_{\tau(d+s)s}(x_{\tau(d+s)1} - x_{\tau(1)1}, \dots, x_{\tau(d+s)d} - x_{\tau(1)d})$. We claim that $\text{Rank}(\mathbf{A}_{\tau d}^{(3)}) \leq d-1$. Otherwise, if $\text{Rank}(\mathbf{A}_{\tau d}^{(3)}) = d$, then there exist $i_1, \dots, i_d \in \{2, \dots, d+J-1\}$, such that, the sub-matrix consisting of the i_1 th, \dots , i_d th rows of $\mathbf{A}_{\tau d}^{(3)}$ is nonsingular. Then the sub-matrix consisting of the $\tau(1)$ th, $\tau(i_1)$ th, \dots , $\tau(i_d)$ th rows of \mathbf{X}_1 is nonsingular, which implies $\text{Rank}(\mathbf{X}_1) = d+1$. The contradiction implies $\text{Rank}(\mathbf{A}_{\tau d}^{(3)}) \leq d-1$. Then $|\mathbf{A}_\tau^{(3)}| = 0$ and thus $|\mathbf{A}_\tau| = 0$ for each τ . Based on Theorem 2, $|\mathbf{F}| \equiv 0$ and thus $f(\mathbf{p}) \equiv 0$. \square

Proof of Theorem 5 Combining Theorem 1 and Theorem 4, it is straightforward that $f(\mathbf{p}) = 0$ if $\text{Rank}(\mathbf{X}_1[\{i \mid p_i > 0\}]) \leq d$. We only need to show that $f(\mathbf{p}) > 0$ if $\text{Rank}(\mathbf{X}_1[\{i \mid p_i > 0\}]) = d+1$. Due to Theorem 1, we only need to verify the case $p_i > 0, i = 1, \dots, m$, because otherwise we may simply remove all support points with $p_i = 0$.

Suppose $p_i > 0, i = 1, \dots, m$ and $\text{Rank}(\mathbf{X}_1) = d+1$. Then there exist $i_0, \dots, i_d \in \{1, \dots, m\}$, such that, $|\mathbf{X}_1[i_0, \dots, i_d]| \neq 0$. According to the proof of Theorem 4, for each $i \in \{i_0, \dots, i_d\}$, there exists an $\epsilon_i \in (0, 1)$, such that, $f(\mathbf{p}) > 0$ as long as $p_i = x \in (1 - \epsilon_i, 1)$ and $p_j = (1 - x)/(m - 1)$ for $j \neq i$. On the other hand, for each $i \notin \{i_0, \dots, i_d\}$, if we denote the j th row of \mathbf{X}_1 by $\alpha_j, j = 1, \dots, m$, then $\alpha_i = a_0\alpha_{i_0} + \dots + a_d\alpha_{i_d}$ for some real numbers a_0, \dots, a_d . Since $\alpha_i \neq 0$, then at least one $a_i \neq 0$. Without any loss of generality, we assume $a_0 \neq 0$. Then it can be verified that $|\mathbf{X}_1[i, i_1, \dots, i_d]| \neq 0$ too. Following the proof of Theorem 4 again, for such an $i \notin \{i_0, \dots, i_d\}$, there also exists an $\epsilon_i \in (0, 1)$, such that, $f(\mathbf{p}) > 0$ as long as $p_i = x \in (1 - \epsilon_i, 1)$ and $p_j = (1 - x)/(m - 1)$ for $j \neq i$. Let $\epsilon_* = \min\{\min_i \epsilon_i, (m - 1) \min_i p_i, 1 - 1/m\}/2$. For $i = 1, \dots, m$, denote $\delta_i = (\delta_{i1}, \dots, \delta_{im})^T \in S$ with $\delta_{ii} = 1 - \epsilon_*$ and $\delta_{ij} = \epsilon_*/(m - 1)$ for $j \neq i$. It can be verified that $\mathbf{p} = a_1\delta_1 + \dots + a_m\delta_m$ with $a_i = (p_i - \epsilon_*/(m - 1))/(1 -$

$m\epsilon_*/(m-1)$). By the choice of ϵ_* , $f(\delta_i) > 0$, $a_i > 0$, $i = 1, \dots, m$, and $\sum_i a_i = 1$. Then $f(\mathbf{p}) > 0$ according to Lemma S.6. \square

Proof of Corollary 3 In order to check when a minimally supported design supported only on $\{x_1, x_2\}$ is D-optimal, we add one more support point, that is, x_3 . According to Theorem 2, Lemmas S.4, S.5, and 2, the objective function for a D-optimal approximate design on $\{x_1, x_2, x_3\}$ is $f(p_1, p_2, p_3) = p_1 p_2 (c_{210} p_1 + c_{120} p_2) + p_1 p_3 (c_{201} p_1 + c_{102} p_3) + p_2 p_3 (c_{021} p_2 + c_{012} p_3) + c_{111} p_1 p_2 p_3$, where

$$\begin{aligned} c_{210} &= e_2 g_{11}^2 g_{12}^2 (\pi_{11} \pi_{12} \pi_{13})^{-1} (x_1 - x_2)^2 > 0 \\ c_{120} &= e_1 g_{21}^2 g_{22}^2 (\pi_{21} \pi_{22} \pi_{23})^{-1} (x_1 - x_2)^2 > 0 \\ c_{201} &= e_3 g_{11}^2 g_{12}^2 (\pi_{11} \pi_{12} \pi_{13})^{-1} (x_1 - x_3)^2 > 0 \\ c_{102} &= e_1 g_{31}^2 g_{32}^2 (\pi_{31} \pi_{32} \pi_{33})^{-1} (x_1 - x_3)^2 > 0 \\ c_{021} &= e_3 g_{21}^2 g_{22}^2 (\pi_{21} \pi_{22} \pi_{23})^{-1} (x_2 - x_3)^2 > 0 \\ c_{012} &= e_2 g_{31}^2 g_{32}^2 (\pi_{31} \pi_{32} \pi_{33})^{-1} (x_2 - x_3)^2 > 0 \\ c_{111} &= e_1 (u_{22} u_{31} + u_{21} u_{32} - 2b_{22} b_{32}) (x_1 - x_2) (x_1 - x_3) + \\ &\quad e_2 (u_{12} u_{31} + u_{11} u_{32} - 2b_{12} b_{32}) (x_2 - x_1) (x_2 - x_3) + \\ &\quad e_3 (u_{12} u_{21} + u_{11} u_{22} - 2b_{12} b_{22}) (x_3 - x_1) (x_3 - x_2) \end{aligned}$$

Based on Theorem 10, the design $\mathbf{p} = (p_1^*, p_2^*, 0)^T$ is D-optimal if and only if

$$\partial f(\mathbf{p}) / \partial f(p_1) = \partial f(\mathbf{p}) / \partial f(p_2) \geq \partial f(\mathbf{p}) / \partial f(p_3)$$

Similar conclusions could be justified for x_4, \dots, x_m if $m \geq 4$. \square

Proof of Theorem 12 According to the solutions provided by the software **Mathematica**, the largest root of equation (S.10) after simplification is

$$y_1 = -\frac{b_2}{3} - \frac{2^{1/3}(3b_1 - b_2^2)}{3A^{1/3}} + \frac{A^{1/3}}{3 \times 2^{1/3}} \quad (\text{S.11})$$

where $A = -27b_0 + 9b_1 b_2 - 2b_2^3 + 3^{3/2}(27b_0^2 + 4b_1^3 - 18b_0 b_1 b_2 - b_1^2 b_2^2 + 4b_0 b_2^3)^{1/2}$, and $b_i = c_i/c_3$, $i = 0, 1, 2$. Note that the calculation of A and thus y_1 should be regarded as operations among complex numbers since the expression under square root could be negative. Nevertheless, y_1 at the end would be a real number. Thus we are able to provide the analytic solution maximizing $f(p_1, p_2, p_3)$. \square

Proof of Corollary 5 In order to check when a minimally supported design is D-optimal, we first add the four design points, that is, we consider

four design points (x_{i1}, x_{i2}) , $i = 1, 2, 3, 4$ and check when the D-optimal design could be constructed on the first three design points. Let \mathbf{X}_1 be defined as in Lemma 2. In this case, \mathbf{X}_1 is a 4×3 matrix. Following Theorem 2, Lemmas S.4, S.5, and 2, the objective function for a minimally supported design at $(d, J, m) = (2, 3, 4)$ is

$$\begin{aligned} f(p_1, p_2, p_3, p_4) &= c_{1111}p_1p_2p_3p_4 \\ &+ |\mathbf{X}_1[1, 2, 3]|^2 e_1e_2e_3 \cdot p_1p_2p_3(w_1p_1 + w_2p_2 + w_3p_3) \\ &+ |\mathbf{X}_1[1, 2, 4]|^2 e_1e_2e_4 \cdot p_1p_2p_4(w_1p_1 + w_2p_2 + w_4p_4) \\ &+ |\mathbf{X}_1[1, 3, 4]|^2 e_1e_3e_4 \cdot p_1p_3p_4(w_1p_1 + w_3p_3 + w_4p_4) \\ &+ |\mathbf{X}_1[2, 3, 4]|^2 e_2e_3e_4 \cdot p_2p_3p_4(w_2p_2 + w_3p_3 + w_4p_4) \end{aligned}$$

where $e_i = u_{i1} + u_{i2} - 2b_{i2}$, $w_i = e_i^{-1}g_{i1}^2g_{i2}^2(\pi_{i1}\pi_{i2}\pi_{i3})^{-1}$, $i = 1, 2, 3, 4$, and

$$c_{1111} = \sum_{1 \leq i < j \leq 4} e_i e_j (u_{k1}u_{l2} + u_{k2}u_{l1} - 2b_{k2}b_{l2}) \cdot |\mathbf{X}_1[i, j, k]| \cdot |\mathbf{X}_1[i, j, l]| \quad (\text{S.12})$$

with $\{i, j, k, l\} = \{1, 2, 3, 4\}$ given $1 \leq i < j \leq 4$.

According to Theorem 10, a minimally supported design $\mathbf{p} = (p_1^*, p_2^*, p_3^*, 0)^T$ in this case is D-optimal if and only if $\partial f/\partial p_1 = \partial f/\partial p_2 = \partial f/\partial p_3 \geq \partial f/\partial p_4$ at \mathbf{p} . Then $\partial f/\partial p_1 = \partial f/\partial p_2 = \partial f/\partial p_3$ at \mathbf{p} is equivalent to (1) of Corollary 5, and $\partial f/\partial p_4 \leq \partial f/\partial p_1$ at \mathbf{p} leads to (2) of Corollary 5 since the forms of $\partial f/\partial p_i$ at \mathbf{p} , $i = 1, 2, 3$ will not change if more than four design points (i.e., $m > 4$) are added into consideration. Note that

$|\mathbf{X}_1[1, 2, 3]|^2 e_1e_2e_3p_2^*p_3^*(2w_1p_1^* + w_2p_2^* + w_3p_3^*)$ in (2) of Corollary 5 is equal to $\partial f/\partial p_1$ at \mathbf{p} . It could be replaced with $|\mathbf{X}_1[1, 2, 3]|^2 e_1e_2e_3p_1^*p_3^*(w_1p_1^* + 2w_2p_2^* + w_3p_3^*)$ (i.e., $\partial f/\partial p_2$), or $|\mathbf{X}_1[1, 2, 3]|^2 e_1e_2e_3p_1^*p_2^*(w_1p_1^* + w_2p_2^* + 2w_3p_3^*)$ (i.e., $\partial f/\partial p_3$), since these three are all equal. \square

S.4 Maximization of $f_i(z)$ in Section 3

According to Theorem 6, $f_i(z)$ is an order- $(d + J - 1)$ polynomial of z . In order to determine its coefficients a_0, a_1, \dots, a_{J-1} as in (3.2), we need to calculate $f_i(0), f_i(1/2), f_i(1/3), \dots, f_i(1/J)$, which are J determinants defined in (3.1).

Note that \mathbf{B}_{J-1}^{-1} is a matrix determined by $J - 1$ only. For example, $B_1^{-1} = 1$ for $J = 2$,

$$B_2^{-1} = \begin{pmatrix} 2 & -1 \\ -1 & 1 \end{pmatrix}, B_3^{-1} = \begin{pmatrix} 3 & -3 & 1 \\ -\frac{5}{2} & 4 & -\frac{3}{2} \\ \frac{1}{2} & -1 & \frac{1}{2} \end{pmatrix},$$

$$B_4^{-1} = \begin{pmatrix} 4 & -6 & 4 & -1 \\ -\frac{13}{3} & \frac{19}{2} & -7 & \frac{11}{6} \\ \frac{3}{2} & -4 & \frac{7}{2} & -1 \\ -\frac{1}{6} & \frac{1}{2} & -\frac{1}{2} & \frac{1}{6} \end{pmatrix}$$

for $J = 3, 4$, or 5 respectively.

Once a_0, \dots, a_{J-1} in (3.2) are determined, the maximization of $f_i(z)$ on $z \in [0, 1]$ is numerically straightforward since it is a polynomial and its derivative $f'_i(z)$ is given by

$$(1-z)^d \sum_{j=1}^{J-1} j a_j z^{j-1} (1-z)^{J-1-j} - (1-z)^{d-1} \sum_{j=0}^{J-1} (d+J-1-j) a_j z^j (1-z)^{J-1-j} \quad (\text{S.13})$$

S.5 Exchange algorithm for D-optimal exact allocation in Section 4

Exchange algorithm for D-optimal allocation $(n_1, \dots, n_m)^T$ given $n > 0$:

- 1° Start with an initial design $\mathbf{n} = (n_1, \dots, n_m)^T$ such that $f(\mathbf{n}) > 0$.
- 2° Set up a random order of (i, j) going through all pairs $\{(1, 2), (1, 3), \dots, (1, m), (2, 3), \dots, (m-1, m)\}$.
- 3° For each (i, j) , let $c = n_i + n_j$. If $c = 0$, let $\mathbf{n}_{ij}^* = \mathbf{n}$. Otherwise, there are two cases. *Case one:* $0 < c \leq J$, we calculate $f_{ij}(z)$ as defined in (4.1) for $z = 0, 1, \dots, c$ directly and find z^* which maximizes $f_{ij}(z)$. *Case two:* $c > J$, we first calculate $f_{ij}(z)$ for $z = 0, 1, \dots, J$; secondly determine c_0, c_1, \dots, c_J in (4.2) according to Theorem 9; thirdly calculate $f_{ij}(z)$ for $z = J+1, \dots, c$ based on (4.2); fourthly find z^* maximizing $f_{ij}(z)$ for $z = 0, \dots, c$. For both cases, we define

$$\mathbf{n}_{ij}^* = (n_1, \dots, n_{i-1}, z^*, n_{i+1}, \dots, n_{j-1}, c - z^*, n_{j+1}, \dots, n_m)^T$$

Note that $f(\mathbf{n}_{ij}^*) = f_{ij}(z^*) \geq f(\mathbf{n}) > 0$. If $f(\mathbf{n}_{ij}^*) > f(\mathbf{n})$, replace \mathbf{n} with \mathbf{n}_{ij}^* , and $f(\mathbf{n})$ with $f(\mathbf{n}_{ij}^*)$.

- 4° Repeat 2° \sim 3° until convergence, that is, $f(\mathbf{n}_{ij}^*) = f(\mathbf{n})$ in step 3° for any (i, j) .

S.6 More on Example 6: Polysilicon Deposition Study

Table S.1 shows the list of experimental settings for the polysilicon deposition study. The factors are decomposition temperature(A), decomposition pressure(B), nitrogen flow (C), silane flow(D), setting time(E), cleaning method(F). Column 1 provides original indices of experimental settings out of 729 distinct ones. For each experimental setting labelled “1” in a design, 9 responses are collected (Phadke (1989)) and assumed to be independent.

Table S.1: Polysilicon Deposition Study: Experimental Settings for the Original, Rounded Approximate, and D-optimal Exact Designs

Index	A	B	C	D	E	F	Original	Rounded	D-optimal
1	1	1	1	1	1	1	1	0	0
76	1	1	3	3	2	1	1	0	0
89	1	2	1	1	3	2	1	0	0
98	1	2	1	2	3	2	0	0	1
111	1	2	2	1	1	3	0	0	1
116	1	2	2	1	3	2	0	1	0
122	1	2	2	2	2	2	1	0	0
130	1	2	2	3	2	1	0	0	1
167	1	3	1	1	2	2	0	0	1
181	1	3	1	3	1	1	0	1	0
199	1	3	2	2	1	1	0	1	1
201	1	3	2	2	1	3	1	0	0
243	1	3	3	3	3	3	1	0	1
258	2	1	1	2	2	3	1	0	0
286	2	1	2	2	3	1	0	1	0
290	2	1	2	3	1	2	1	0	0
291	2	1	2	3	1	3	0	1	0
294	2	1	2	3	2	3	0	0	1
299	2	1	3	1	1	2	0	0	1
301	2	1	3	1	2	1	0	1	0
313	2	1	3	2	3	1	0	0	1
331	2	2	1	1	3	1	0	1	1
336	2	2	1	2	1	3	0	1	1
339	2	2	1	2	2	3	0	1	0
350	2	2	1	3	3	2	0	1	0
365	2	2	2	2	2	2	0	0	1
376	2	2	2	3	3	1	1	0	0
384	2	2	3	1	2	3	1	0	0
394	2	2	3	2	3	1	0	1	0
399	2	2	3	3	1	3	0	1	0
407	2	3	1	1	1	2	0	0	1
421	2	3	1	2	3	1	1	0	0
461	2	3	3	1	1	2	1	1	0
464	2	3	3	1	2	2	0	1	0
495	3	1	1	1	3	3	0	1	0
501	3	1	1	2	2	3	0	0	1
505	3	1	1	3	1	1	0	0	1
521	3	1	2	1	3	2	0	0	1
522	3	1	2	1	3	3	1	0	0
536	3	1	2	3	2	2	0	1	0
557	3	1	3	2	3	2	1	0	0
558	3	1	3	2	3	3	0	1	0
569	3	2	1	1	1	2	0	1	0
588	3	2	1	3	1	3	1	0	0
625	3	2	3	1	2	1	0	0	1
631	3	2	3	2	1	1	1	0	0
641	3	2	3	3	1	2	0	0	1
671	3	3	1	3	2	2	1	0	0
679	3	3	2	1	2	1	1	0	0

Table S.2 shows the model matrix for the D-optimal design \mathbf{n}_o found for the polysilicon deposition study. In this table, each 3-level factor is represented by its linear component and quadratic component. Thus there are level combinations of 12 predictors.

Table S.2: Polysilicon Deposition Study: Model Matrix for the D-optimal Design

Index	A_1	A_2	B_1	B_2	C_1	C_2	D_1	D_2	E_1	E_2	F_1	F_2
98	-1	1	0	-2	-1	1	0	-2	1	1	0	-2
111	-1	1	0	-2	0	-2	-1	1	-1	1	1	1
130	-1	1	0	-2	0	-2	1	1	0	-2	-1	1
167	-1	1	1	1	-1	1	-1	1	0	-2	0	-2
199	-1	1	1	1	0	-2	0	-2	-1	1	-1	1
243	-1	1	1	1	1	1	1	1	1	1	1	1
294	0	-2	-1	1	0	-2	1	1	0	-2	1	1
299	0	-2	-1	1	1	1	-1	1	-1	1	0	-2
313	0	-2	-1	1	1	1	0	-2	1	1	-1	1
331	0	-2	0	-2	-1	1	-1	1	1	1	-1	1
336	0	-2	0	-2	-1	1	0	-2	-1	1	1	1
365	0	-2	0	-2	0	-2	0	-2	0	-2	0	-2
407	0	-2	1	1	-1	1	-1	1	-1	1	0	-2
501	1	1	-1	1	-1	1	0	-2	0	-2	1	1
505	1	1	-1	1	-1	1	1	1	-1	1	-1	1
521	1	1	-1	1	0	-2	-1	1	1	1	0	-2
625	1	1	0	-2	1	1	-1	1	0	-2	-1	1
641	1	1	0	-2	1	1	1	1	-1	1	0	-2