

# A (my) brief history of Combinatorial Analysis on the $\lambda$ -Calculus

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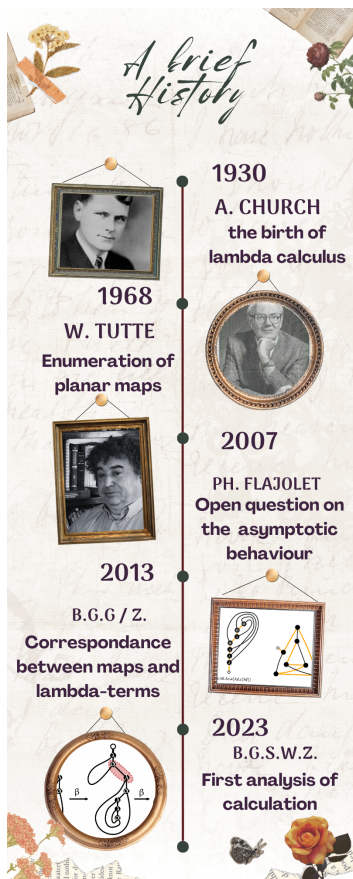
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(Non-exhaustive) list of recent contributors on the topics

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Special thanks to Pr.Dr. Alexandros Singh for many of the figures in this talk

# My plan : A timeline



- What is a  $\lambda$ -terms (for a combinatorist)
- What is a map (for a combinatorist)
- The start of the modern history (for a combinatorist)
- Seems so far away and yet so close
- Towards the grail (for me)

# A. Church and the $\lambda$ -terms

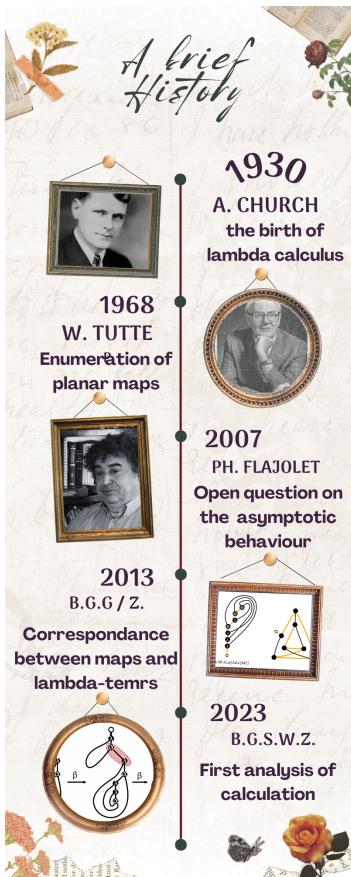
## What is a $\lambda$ -terms ?

$$T ::= a \mid (T * T) \mid \lambda a. T$$

$a$  : variables

$(T * T)$  : application

$\lambda a. T$  : abstraction



# A. Church and the $\lambda$ -terms

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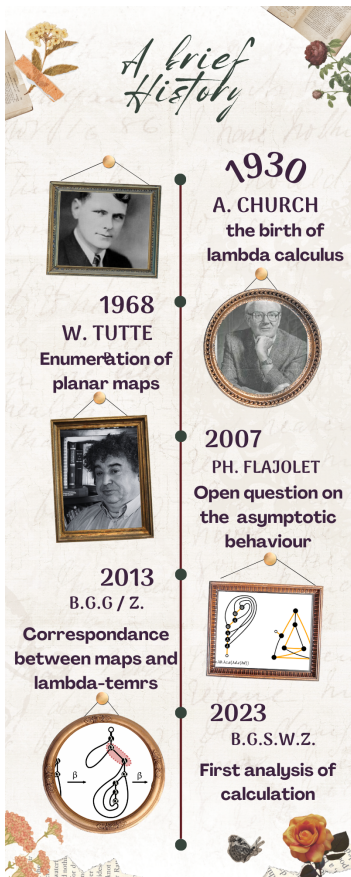
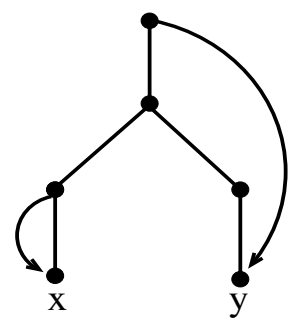
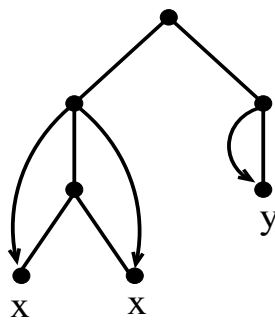
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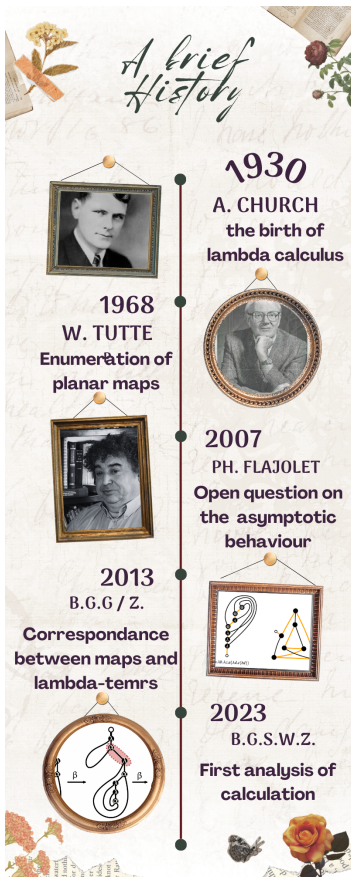
$\lambda a. T$  : abstraction

$$(\lambda x. (x * x) * \lambda y * y)$$

$$\lambda y. (\lambda x. x * \lambda x. y)$$

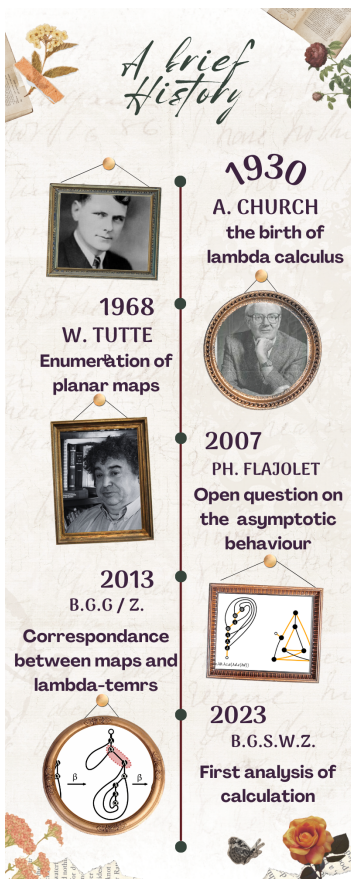


# A. Church and the $\lambda$ -terms



- $\lambda$ -terms can be **close** : no free variable.

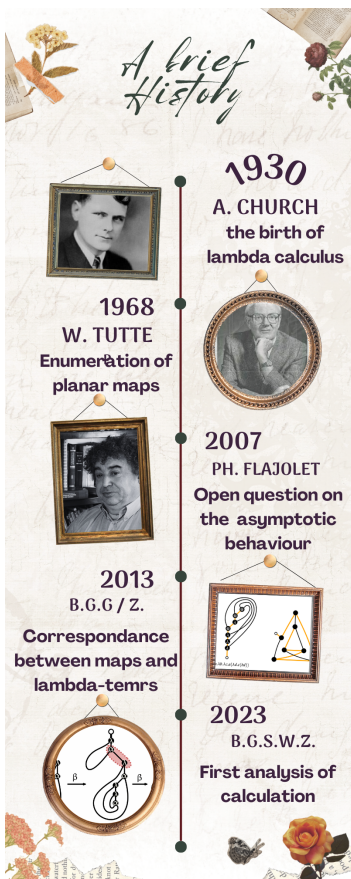
## A. Church and the $\lambda$ -terms



- $\lambda$ -terms can be **close** : no free variable.
- Considered up to  **$\alpha$ -conversion** : renaming of variables.

$$(\lambda x.(x * y)) = (\lambda z.(z * t)) \neq (\lambda y.(y * y))$$

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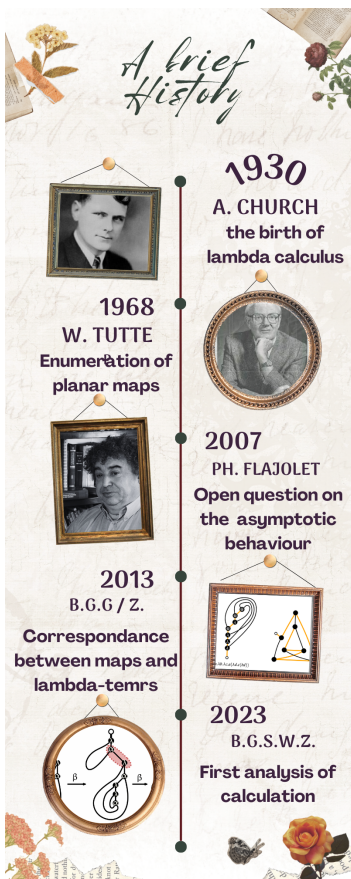
$$(\lambda x.(x * y)) = (\lambda z.(z * t)) \neq (\lambda y.(y * y))$$

- **linear** : each abstraction binds exactly one variable.

$$(\lambda x.x) * (\lambda y.y)$$



## A. Church and the $\lambda$ -terms



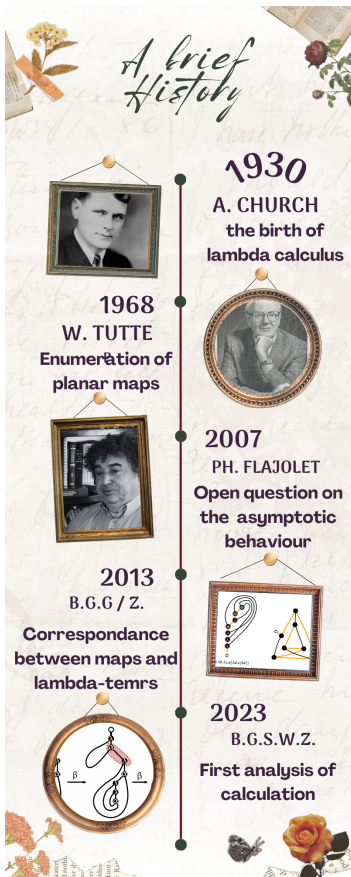
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## A. Church and the $\lambda$ -terms



### Where is the calculation ?

Calculation =  $\beta$ -reduction

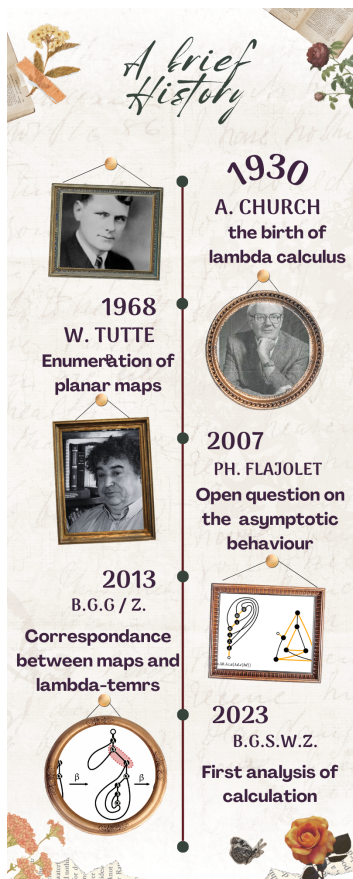
If you have somewhere in your  $\lambda$ -term a **redex** :

$$(\lambda x. T) * Q$$

Then we can apply a  $\beta$ -reduction that corresponds to replace all the instances of  $x$  in  $T$  by  $Q$  :

$$(\lambda x. T) * Q = T[x \leftarrow Q]$$

# A. Church and the $\lambda$ -terms



## Where is the calculation ?

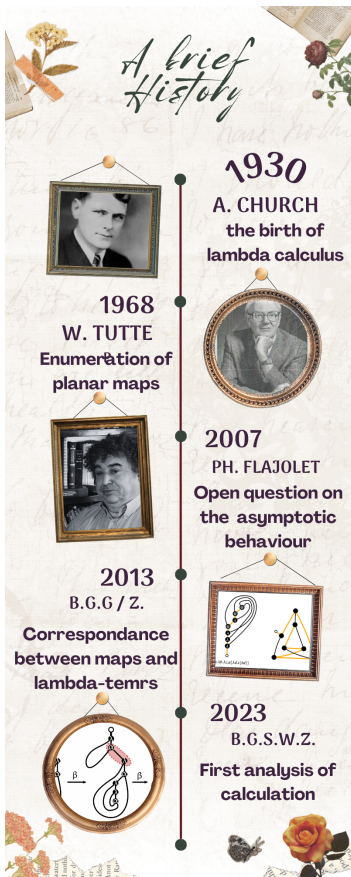
Calculation =  $\beta$ -reduction

Example :

$$S := \lambda t.(\lambda y.\lambda x.(x * x * y)) * (\lambda z.(z * t))$$

Here  $T := x * x * y$  and  $Q := \lambda z.(z * t)$

# A. Church and the $\lambda$ -terms



## Where is the calculation ?

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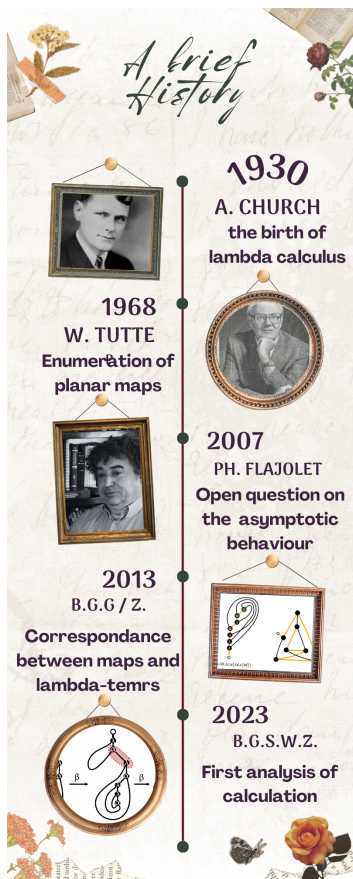
$$S := \lambda t.(\lambda y.\lambda x.(x * x * y)) * (\lambda z.(z * t))$$

Here  $T := x * x * y$  and  $Q := \lambda z.(z * t)$

So after  $\beta$ -reduction, we get

$$S := \lambda t.\lambda y.((\lambda z.(z * t)) * (\lambda z.(z * t)) * y)$$

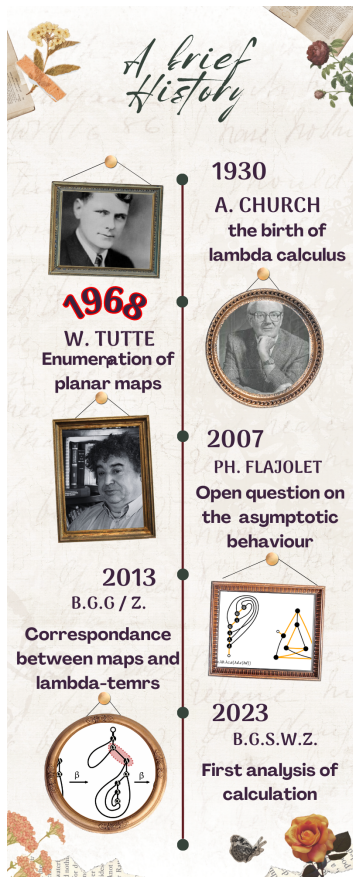
## A. Church and the $\lambda$ -terms



### We have a model of calculation !

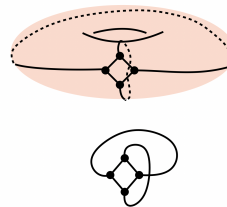
- Base of functional programming by introducing the notion of type.
- Lambda calculation is equivalent in computing power to Turing machines (Turing-Complete).
- Curry-Howard's correspondence between proofs and  $\lambda$ -terms

# W. Tutte and the Maps

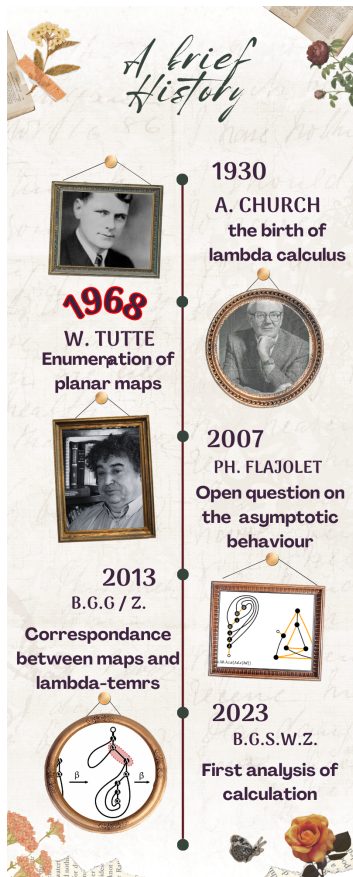


## What is a map ?

A map is the representation of a graph without crossing edges on a surface.

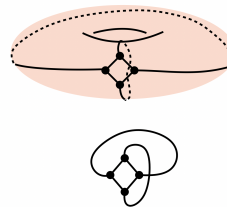


# W. Tutte and the Maps



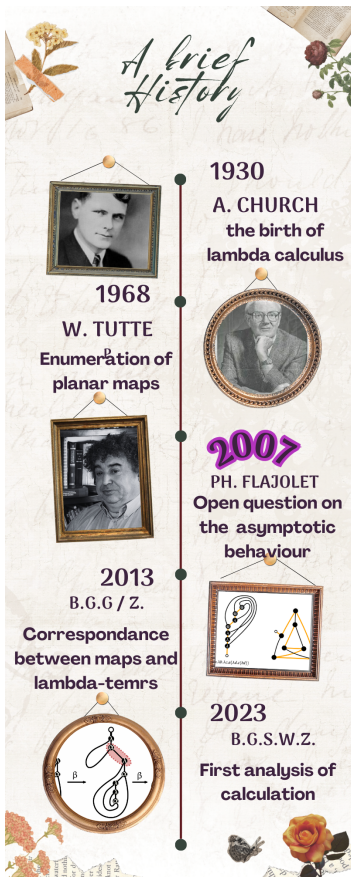
## What is a map ?

A map is the representation of a graph without crossing edges on a surface.



Here, we need maps on any surface (without genus restrictions).

# Ph. Flajolet and the Analytic Combinatoric

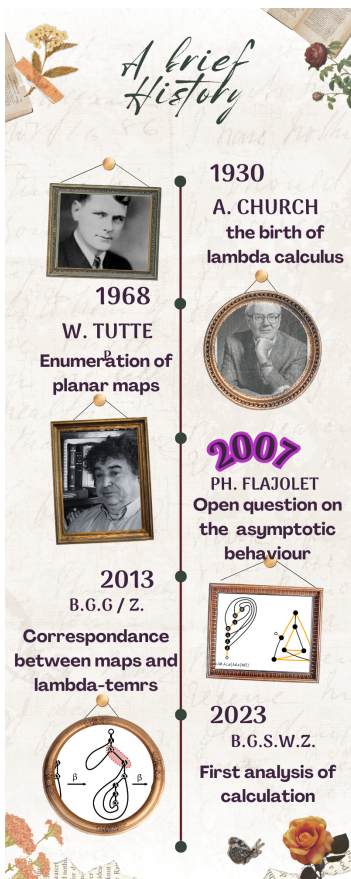


Can we have a combinatorial point of view on  $\lambda$ -calculus ?

- How count the number of  $\lambda$ -terms ?
- What is the asymptotic behaviour ?
- Asymptotic laws of parameters ?



# Ph. Flajolet and the Analytic Combinatoric



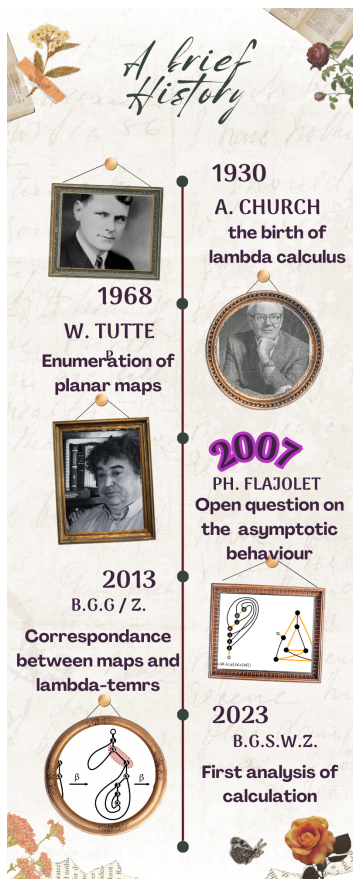
Can we have a combinatorial point of view on  $\lambda$ -calculus ?

Let's try the symbolic method !

$$\mathcal{L} = \mathcal{F} + (\mathcal{N} \times \mathcal{L}^2) + (\mathcal{U} \times \text{subs}(\mathcal{F} \rightarrow \mathcal{F} + \mathcal{B}, \mathcal{L}))$$

- $\mathcal{L}$  : the class of  $\lambda$ -terms with free variables
- $\mathcal{N}$  for the binary nodes (applications)
- $\mathcal{U}$  for the unary nodes (abstractions)
- $\mathcal{F}$  for the free leaves
- $\mathcal{B}$  for the binded leaves

# Ph. Flajolet and the Analytic Combinatoric



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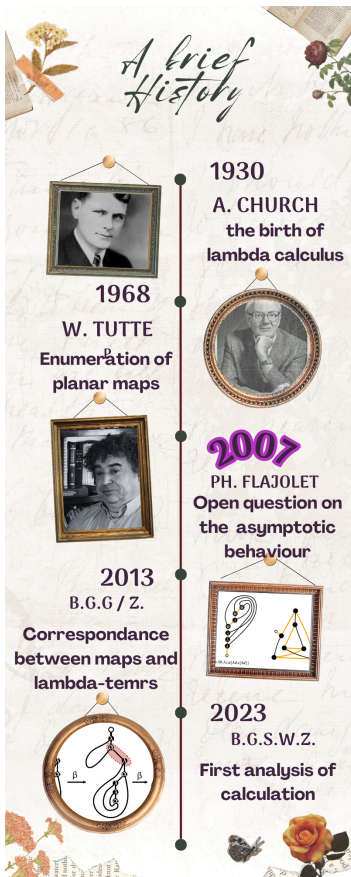
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## ■ Generating function

$$L(z, f) = fz + zL(z, f)^2 + zL(z, f + 1).$$

(size = number of nodes)

# Ph. Flajolet and the Analytic Combinatoric



Can we have a combinatorial point of view on  $\lambda$ -calculus ?

$$L(z, f) = fz + zL(z, f)^2 + zL(z, f + 1).$$

So, the series for the closed  $\lambda$ -terms begins as

$$\begin{aligned} L(z, 0) &= [f^0]L(z, f) \\ &= z^2 + 2z^3 + 4z^4 + 13z^5 + 42z^6 + 139z^7 \\ &\quad + 506z^8 + 1915z^9 + 7558z^{10} + \dots \end{aligned}$$

## Ph. Flajolet and the Analytic Combinatoric

$$L(z, 0) = \frac{1}{2z} \left( 1 - \sqrt{\Lambda(z)} \right)$$

with  $\Lambda(z)$  equal to

$$1 - 2z + 2z \sqrt{1 - 2z - 4z^2} + 2z \sqrt{\dots \sqrt{1 - 2z - 4nz^2} + 2z \sqrt{\dots}}$$

## Ph. Flajolet and the Analytic Combinatoric

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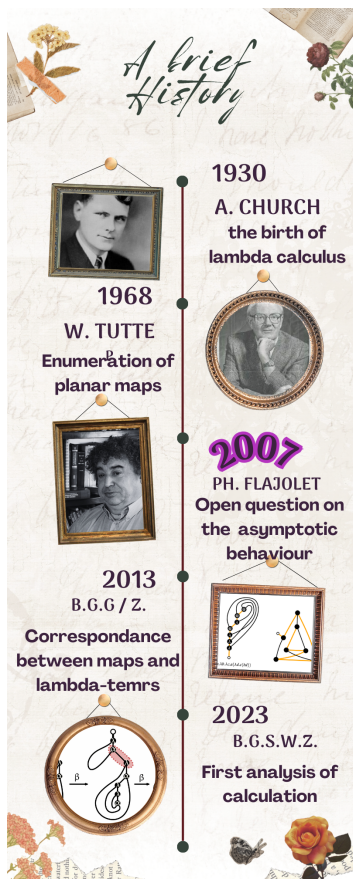
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**First idea** : Restriction on the number of nested radicands...

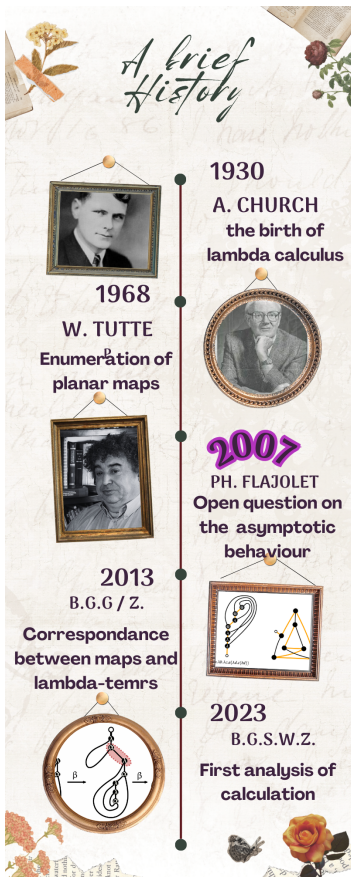
# Ph. Flajolet and the Analytic Combinatoric



## The bounded unary height $\lambda$ -terms

Let fix  $k$  : maximum number of abstractions on a path from the root to a leaf.  
Denote  $S^{(k)}$  the generating function of the  $k$ -bounded unary height  $\lambda$ -terms.

# Ph. Flajolet and the Analytic Combinatoric



## The bounded unary height $\lambda$ -terms

$S^{(k)} = P^{(0,k)}(z)$  is composed of  $k$ -nested radicals (square roots) :

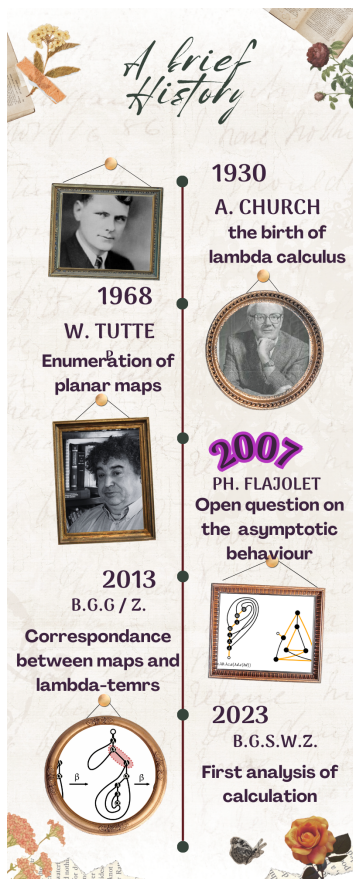
$$P^{(k,k)}(z) = \frac{1 - \sqrt{1 - 4kz^2}}{2z}$$

$$P^{(i,k)}(z) = \frac{1 - \sqrt{1 - 4iz^2 - 4z^2 P^{(i+1,k)}(z)}}{2z}$$

We'll be able to do the asymptotic study of its coefficients !



# Ph. Flajolet and the Analytic Combinatoric



## The bounded unary height $\lambda$ -terms

$S^{(k)}$  is composed of  $k$ -nested radicals :  
Where is "located" the dominant singularity as  $k$  increases ?

- $k = 1$  : innermost radical
- $k = 2$  : second internal radical
- $k = 3, 4, \dots$  : second internal radical
- $k = 9$  : **third** internal radical !

**Damned!!!!**

## Ph. Flajolet and the Analytic Combinatoric

Values of  $k$  for which there are two dominant radicals (coalescence) ?

- Define  $(u_k)_{k \geq 0}$  by  $u_0 = 0$  and

$$u_k = u_{k-1}^2 + k \quad \text{for } k > 0$$

- First values :  $u_1 = 1$ ,  $u_2 = 3$ ,  $u_3 = 12$ ,  $u_4 = 148$ ,  $u_5 = 21909$ , ...
- The sequel  $(u_k)_{k \geq 0}$  is doubly exponential
- $\lim_{k \rightarrow \infty} u_k^{1/2^k} \simeq \chi = 1.36660956\dots$
- $u_k = \lfloor \chi^{2^k} \rfloor$

## Ph. Flajolet and the Analytic Combinatoric

Define :  $N_k = u_k^2 - u_k + k$ .

First values :  $N_1 = 1$ ,  $N_2 = 8$ ,  $N_3 = 135$ ,  $N_4 = 21760$ ,  
 $N_5 = 479982377$ , ...

### Theorem (BGG'2011)

Let  $i$  be such that  $k \in [N_i, N_{i+1}[$ .

By ordering radicals from inner to outer :

- If  $k \neq N_i$ , the dominant radical of  $S^{(k)}(z)$  is the  $i$ -th one ; the dominant singularity is algebraic in type  $1/2$ .
- If  $k = N_i$ , the two radicals of rank  $i$  and  $(i + 1)$  have the same singularity which is dominant ; this dominant singularity is algebraic in type  $1/4$ .

## Ph. Flajolet and the Analytic Combinatoric

## Theorem (BGG'2011)

$[z^n]\mathcal{S}(N_k)_n \sim \frac{1}{\Gamma(3/4)} h_k n^{-5/4} (2u_k)^n$ . when  $n$  tends to  $\infty$

with  $h_k = \left(-\frac{u_k}{2} w_{k-1,k}\right)^{1/4} \prod_{i=k}^{n_k-1} \frac{1}{2u_{-i}}$ .

and  $w_{k-1,k}$  is defined recursively by  $w_{0,k} = -4N_k/u_k$ ,  
 $w_{i,k} = -4(N_k - i)/u_k - 2 + 2u_{k-i}/u_k + w_{i-1,k}/(2u_{k-i})$ .

Numerically :

$$N_1 := 1 : [z^n]\mathcal{S}(1)_n \sim 0.2426128012\dots n^{-5/4} 2^n$$

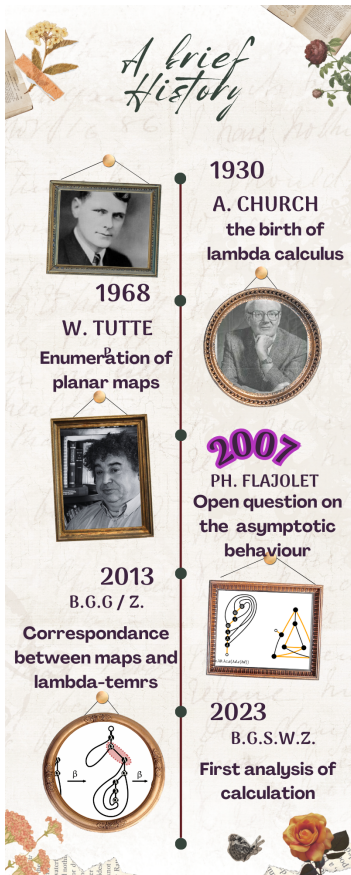
$$N_2 := 8 : [z^n]\mathcal{S}(8)_n \sim 0.00009318885377\dots n^{-5/4} 6^n$$

$$N_3 := 135 : [z^n]\mathcal{S}(135)_n \sim 7.116999389\dots \times 10^{-158} n^{-5/4} 24^n$$

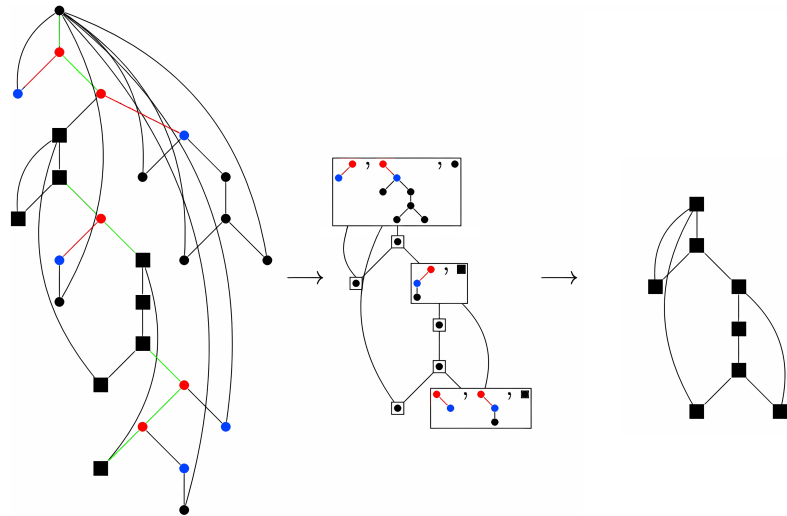
## Ph. Flajolet and the Analytic Combinatoric

Function	Pos. of the dom. radical	Dom. singularity
$S^{(1)}$	$\{1,2\}$	0.5
$S^{(2)}$	2	0.3438
$S^{(3)}$	2	0.2760
...	...	...
$S^{(8)}$	$\{2,3\}$	0.1667
$S^{(9)}$	3	0.1571
...	...	...
$S^{(134)}$	3	0.0418
$S^{(135)}$	$\{3,4\}$	0.0417
$S^{(136)}$	4	0.0415
...	...	...

# Ph. Flajolet and the Analytic Combinatoric

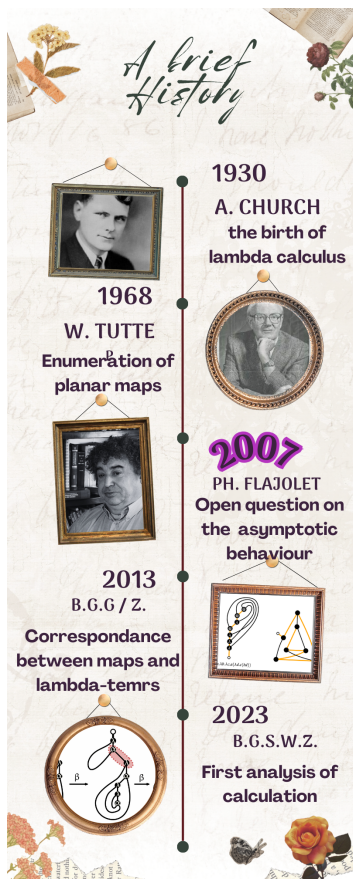


## General $\lambda$ -terms without catalytic variables [BGGJ'13]



$$L(z) = zM(z) + zL(z)^2 + zL\left(\frac{z}{1 - 2zM}\right)$$

# Linear $\lambda$ -terms and maps



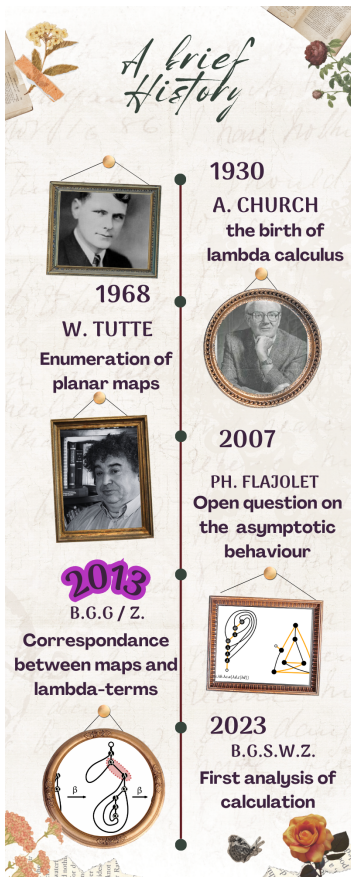
## Asymptotic of the general $\lambda$ -terms ?

$$L(z) = zM(z) + zL(z)^2 + zL\left(\frac{z}{1 - 2zM}\right)$$

$$[z^n]L(z) \asymp \left(\frac{4n}{e \ln(n)}\right)^{n/2}$$

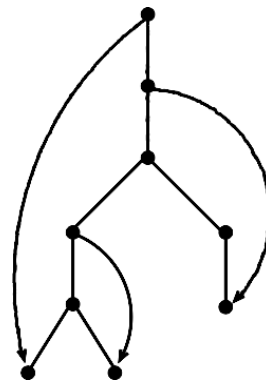
More precise asymptotic is still open

# Linear $\lambda$ -terms and Maps



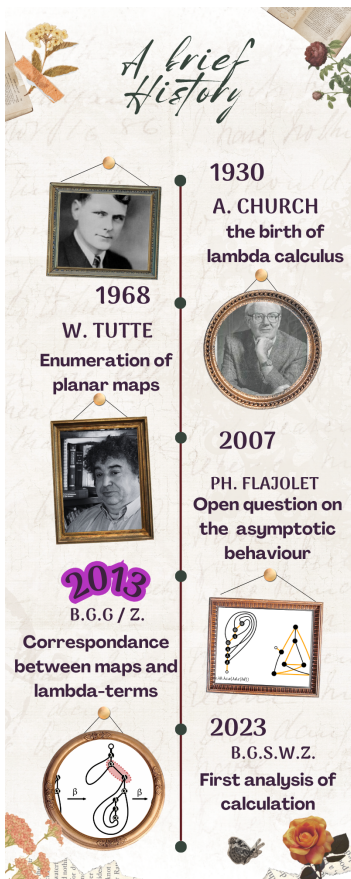
## The sub-class of linear $\lambda$ -terms

Each abstraction binds **exactly** one variable.  
 $n + 1$  variables  $\rightarrow n + 1$  abstractions  $\rightarrow n$  applications  $\rightarrow$  size  $:= 3n + 2$



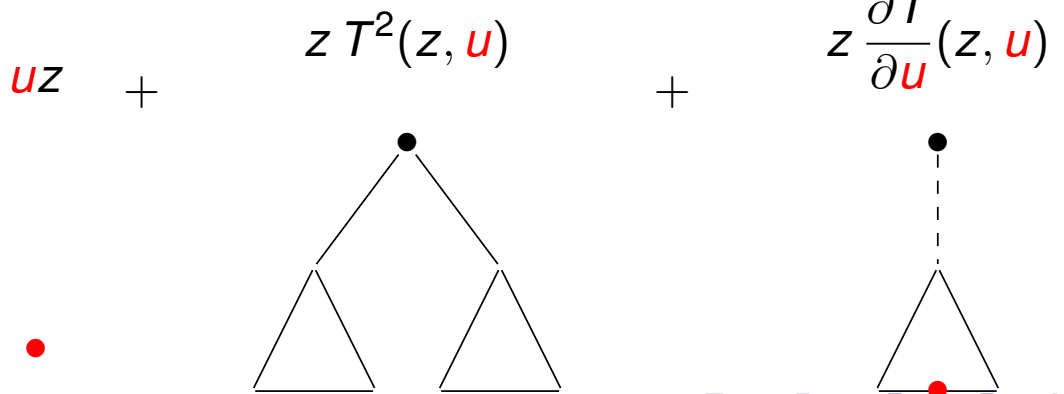


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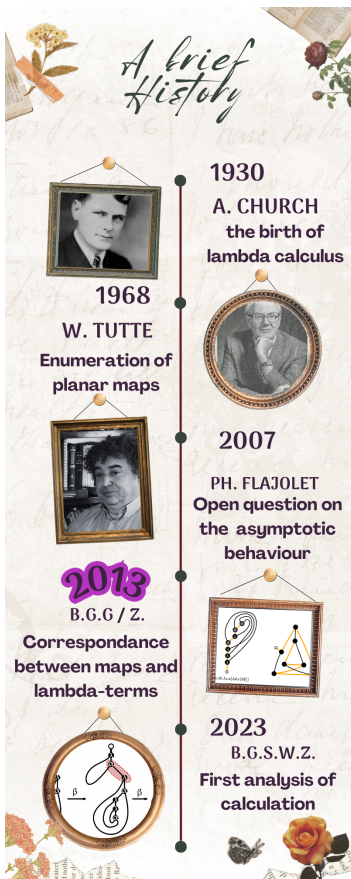


## A catalytic specification for linear $\lambda$ -terms

$$T(z, u) =$$



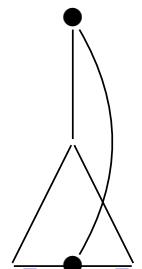
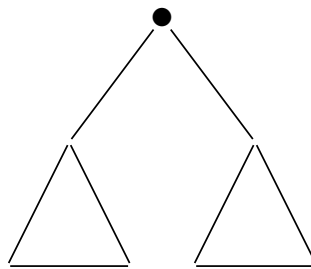
# Linear $\lambda$ -terms and Maps



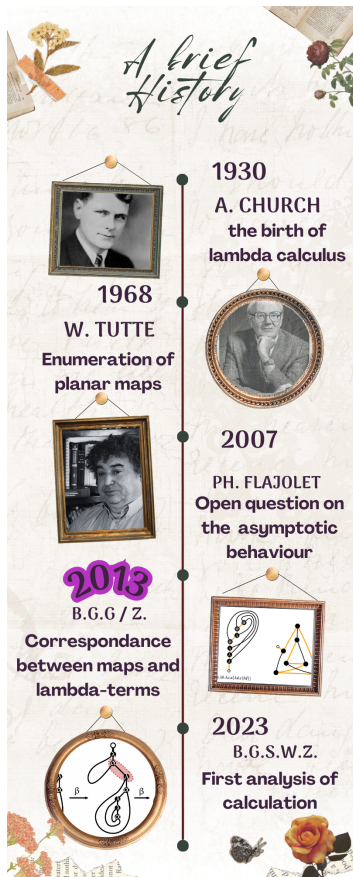
## A catalytic specification for linear $\lambda$ -terms

$$T(z, u) =$$

$$uz + z T^2(z, u) + z \frac{\partial T}{\partial u}(z, u)$$



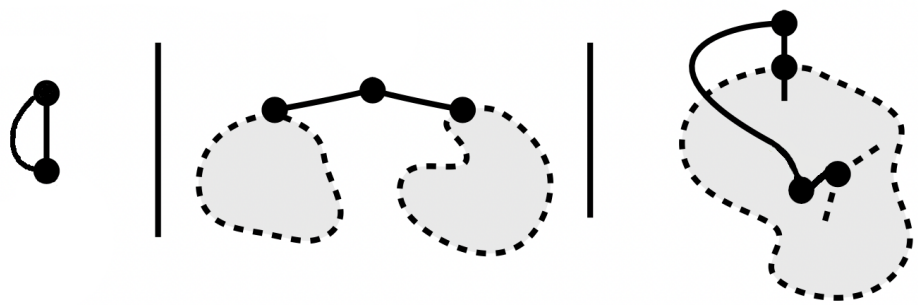
# Linear $\lambda$ -terms and Maps



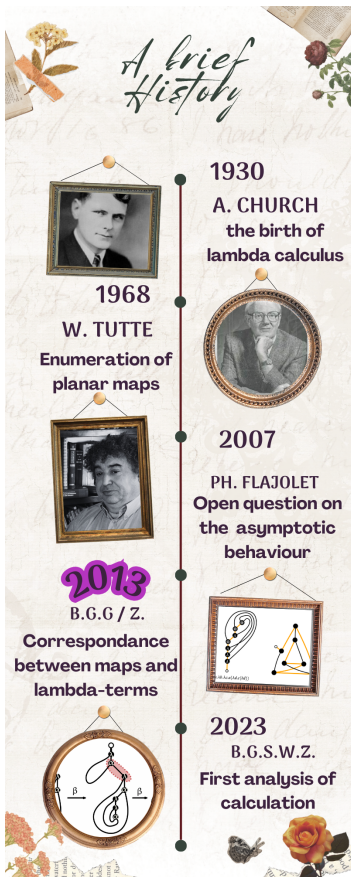
A specification (without catalytic variable) for linear  $\lambda$ -terms

Each abstraction binds **exactly** one variable.

$$T(z) = z^2 + zT(z)^2 + 2z^4 \frac{\partial T}{\partial z}(z)$$



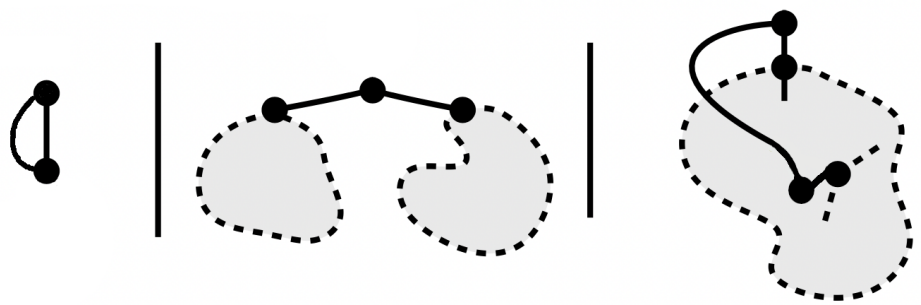
# Linear $\lambda$ -terms and Maps



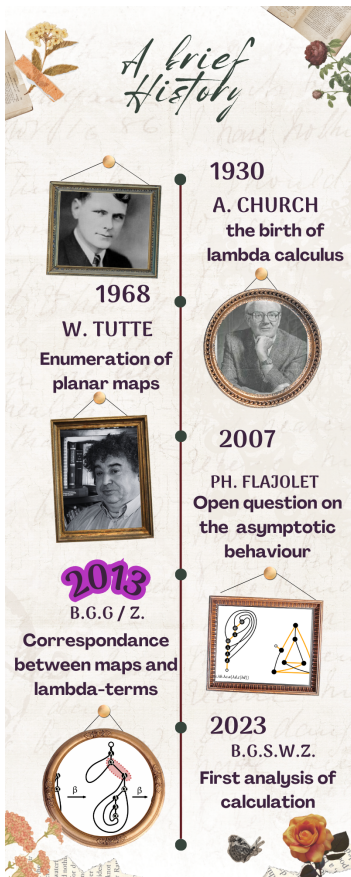
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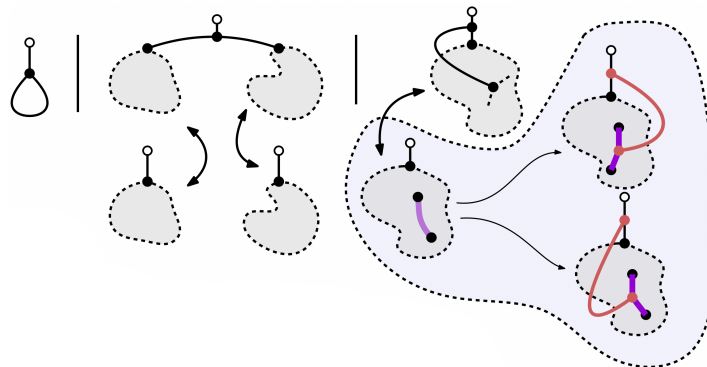


# Linear $\lambda$ -terms and Maps

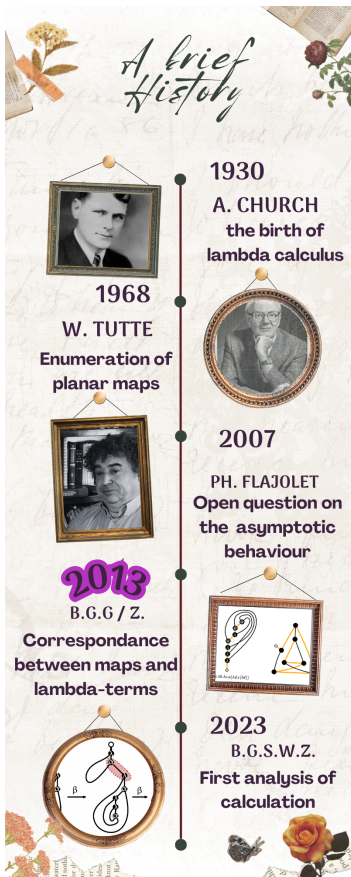


## A specification for connected rooted trivalent maps

$$T(z) = z^2 + zT(z)^2 + 2z^4 \frac{\partial T}{\partial z}(z)$$

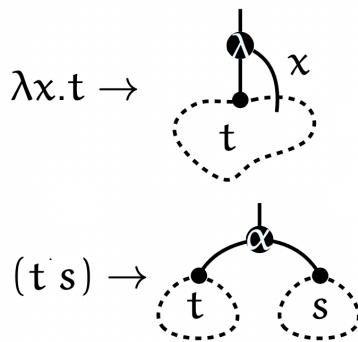


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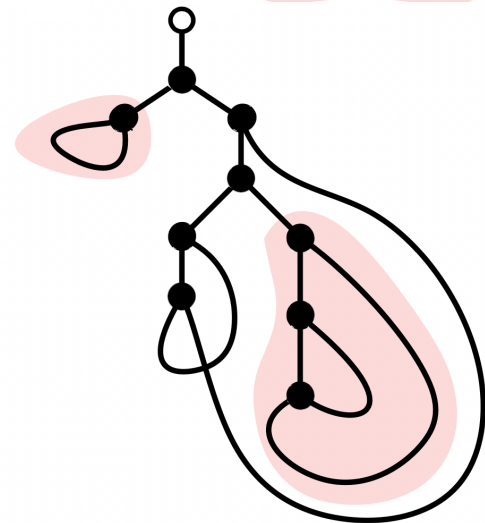


## Linear $\lambda$ -terms $\cong$ rooted trivalent maps

From closed terms to maps



$(\lambda x.x) (\lambda y.(\lambda z.z y) (\lambda w.\lambda u.w u))$



### Dictionary

- # subterms  $\leftrightarrow$  # edges
- closed subterms  $\leftrightarrow$  bridges
- using variables in order  $\leftrightarrow$  planarity of maps

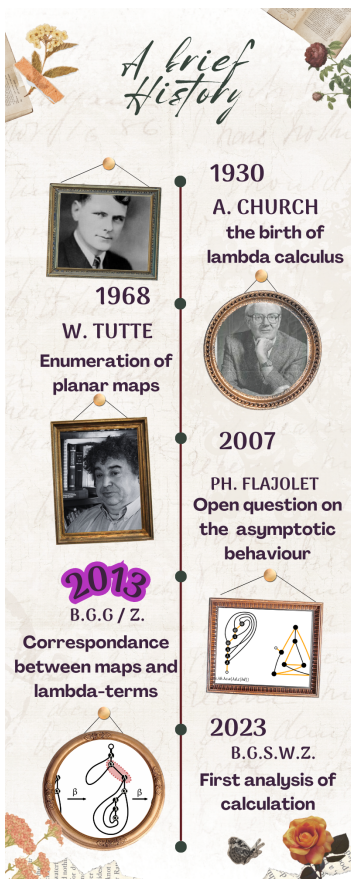
# Divergent Riccati equation

## Divergent Riccati equation

Classical change of variable  $T = \frac{U'}{U}$

Which can be seen as a combinatorial operation :

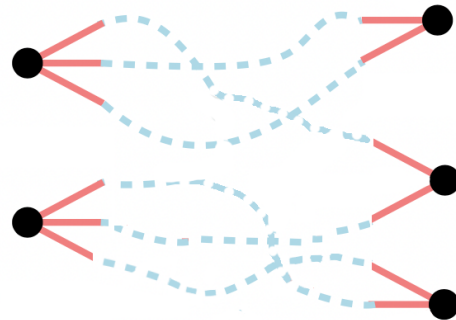
Take only the connected objects + Pointing.  
 All Riccati equation arising in combinatorics (I know) come from that (irreducible diagrams, irreducible permutations,...) !



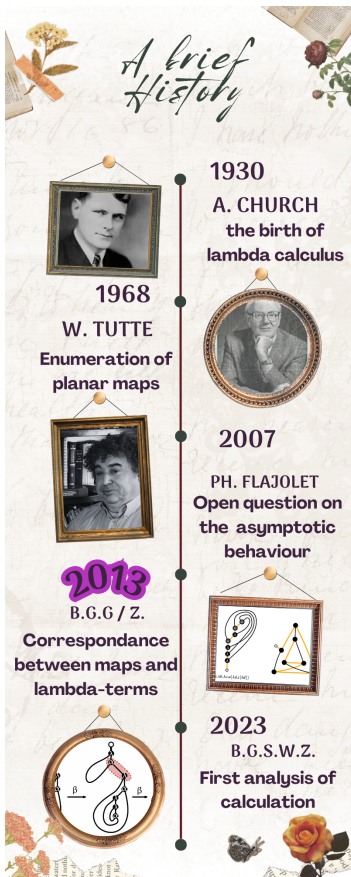
# Divergent Riccati equation

## Divergent Riccati equation

$$T(z^2) = z^4 + z^5 \frac{\partial}{\partial z} \ln \left( \exp(z^3/3) \odot \exp(z^2/2) \right)$$



(rigidity)





# Divergent Riccati equation

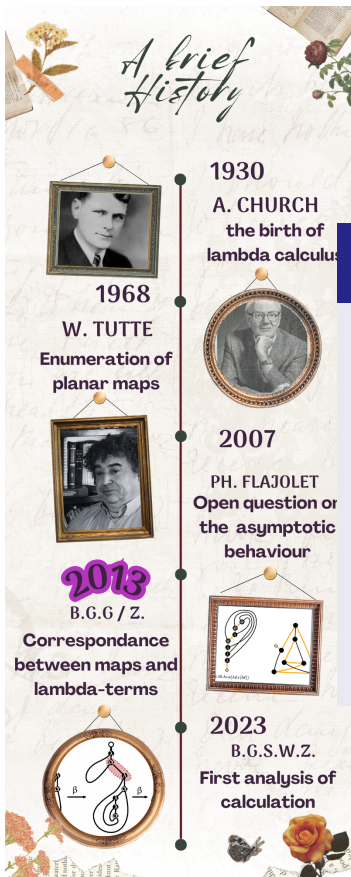
## Divergent Riccati equation

### Theorem (BGJ'2013)

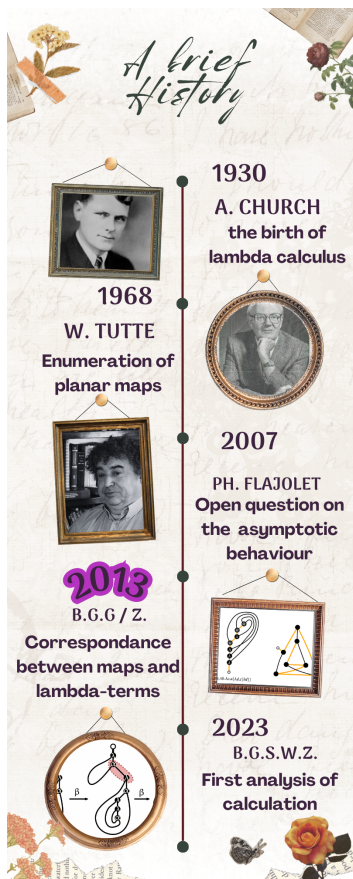
$$T(z^2) = z^4 + z^5 \frac{\partial}{\partial z} \ln \left( \exp(z^3/3) \odot \exp(z^2/2) \right)$$

$$\lambda_{3n+2} \sim 3/\pi \times 6^n n!$$

- Efficient random sampling.

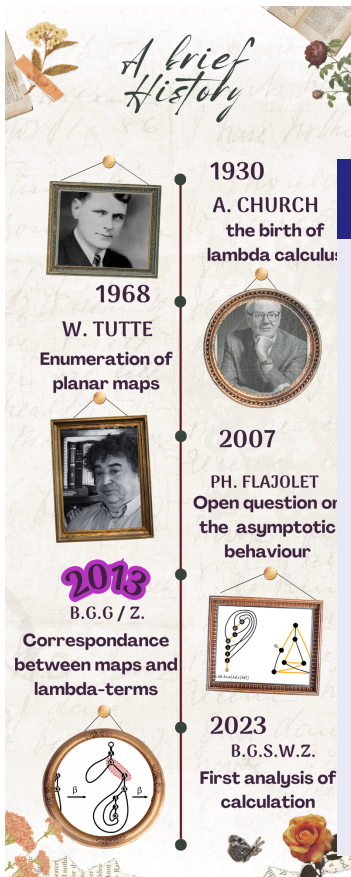


# Linear $\lambda$ -terms and Maps



## Efficient random sampling for linear $\lambda$ -terms

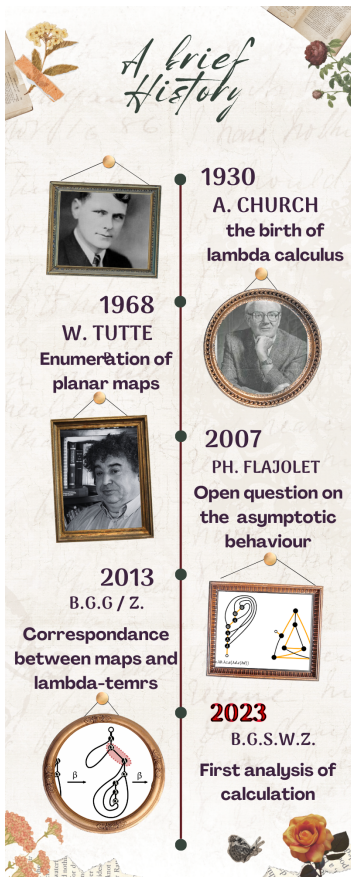
- Do...
- Draw a permutation of size  $6n$  product of 3-cycles.
- Draw an involution without fixed point of size  $6n$ .
- Build a map by gluing half-edges.
- Until the map is connected (this appends with probability 1 when  $n \rightarrow \infty$ )
- Build the linear lambda term from it.

Linear  $\lambda$ -terms and MapsParameters on linear  $\lambda$ -terms

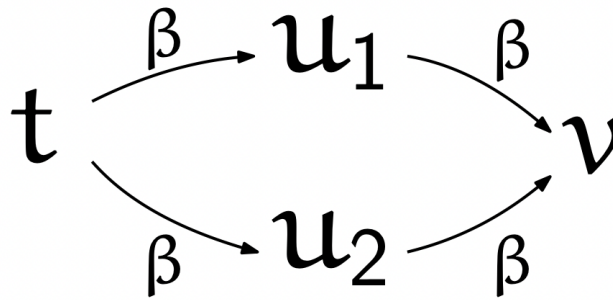
## Theorem (BSZ'2021)

- *The random variable for identity-subterms in closed linear terms converges in distribution to a Poisson of parameter 1.*
- *The random variable for closed subterms in closed linear terms converges in distribution to a Poisson of parameter 1.*
- *The random variable for free variables in open linear terms converges in distribution to  $\mathcal{N}((2n)^{1/3}, (2n)^{1/3})$ .*

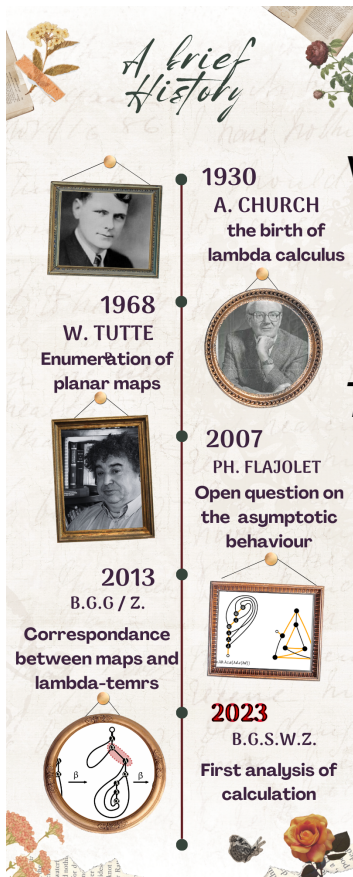
# Towards calculation on linear $\lambda$ -terms



For linear terms,  $\beta$ -reduction possesses the strong diamond property.



All the paths of  $\beta$ -reduction terminates after a finite number of steps to the same  $\lambda$ -term with no redex (called **normal form**). All these paths have same number of steps.

Towards calculation on linear  $\lambda$ -terms

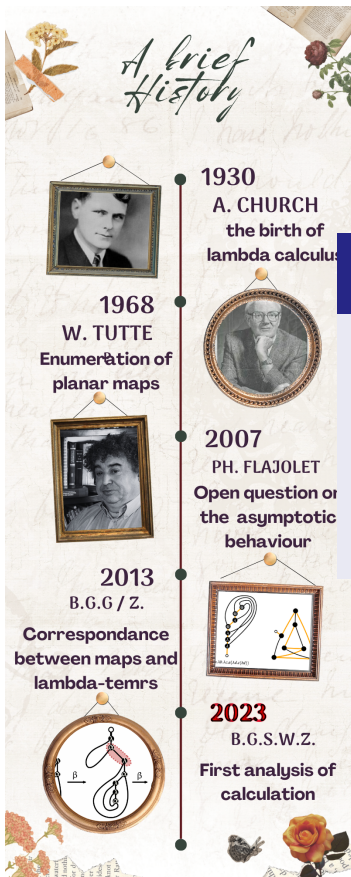
What is the mean length of full  $\beta$ -reduction chain?  
 First step : count (mark with  $r$ ) the number of redices !

$$T = z^2 + zT^2 +$$

$$z^3(1 + (r-1)zT) \left( \frac{z(r+5)\partial_z T}{3} - (r^2-1)\partial_r T \right) \\ + \frac{z^4(r-1)^2 T^2}{3} + \frac{4z^3(r-1)T}{3}$$

Then, pump the first two moments...

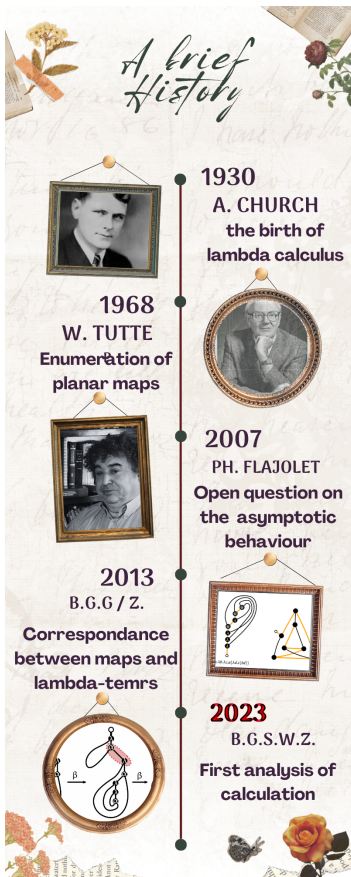
# Towards calculation on linear $\lambda$ -terms



## Theorem (BSZ'2021)

*Mean number of redices in closed linear terms of size  $n$  is asymptotically in  $\frac{n}{24}$*   
*The variance is also asymptotically in  $\frac{n}{24}$*

This is of course a first lower bound for the length of the  $\beta$ -redex chain...

Towards calculation on linear  $\lambda$ -terms

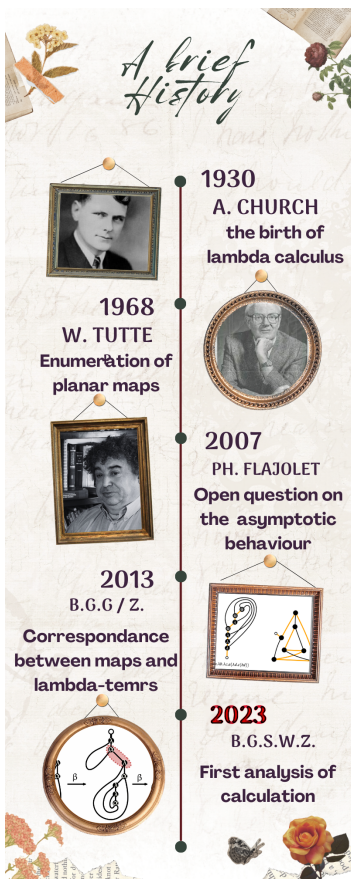
The only three patterns that "create" a new redex :  
(see JJ Lévy's thesis)

$$\rho_1 := (\lambda x. C[(xt_1]))(\lambda y. t_2) \rightarrow_{\beta} C[((\lambda y. t_2)t_1)]$$

$$\rho_2 := (\lambda x. x)(\lambda y. t_1)t_2 \rightarrow_{\beta} (\lambda y. t_1)t_2$$

$$\rho_3 := ((\lambda x. \lambda y. t_1)t_2)t_3 \rightarrow_{\beta} (\lambda y. t_1[x := t_2])t_3$$

# Towards calculation on linear $\lambda$ -terms

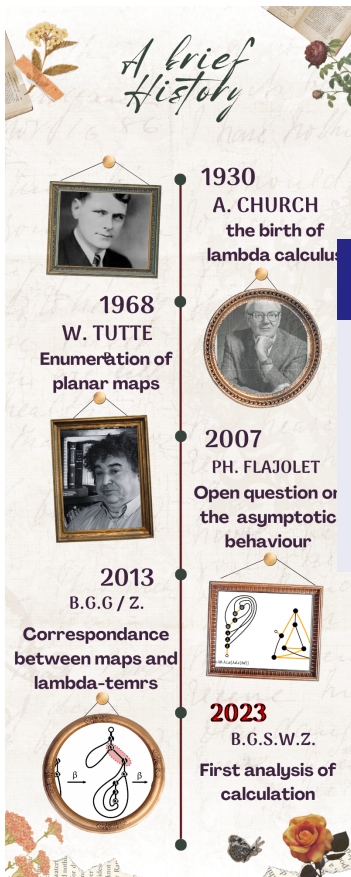


The length of the  $\beta$ -reduction chain is larger than  $|\beta\text{-redex}| + |p_1| + |p_2| + |p_3|$

All that's left to do is work...



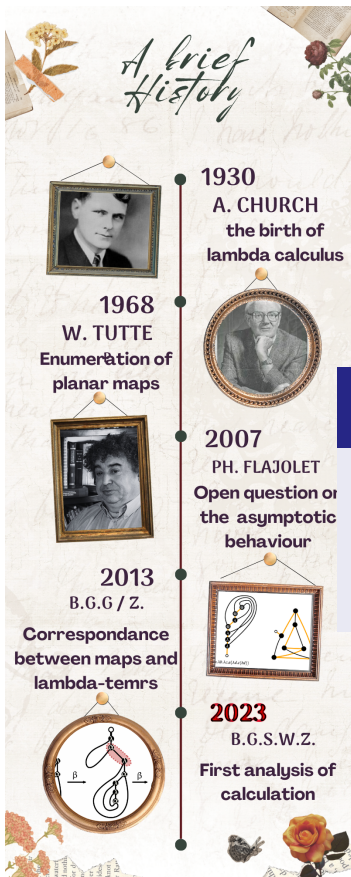
# Towards calculation on linear $\lambda$ -terms



## Theorem (BGSWZ'2023)

The mean number of pattern  $p_1$  (resp.  $p_2$ ) in closed linear terms of size  $n$  is asymptotically a constant  $\mu_1 = \frac{1}{6}$  (reps.  $\mu_2 = \frac{1}{48}$ )

Essentially, we still are able to specify these patterns (but that needs some care !)

Towards calculation on linear  $\lambda$ -terms

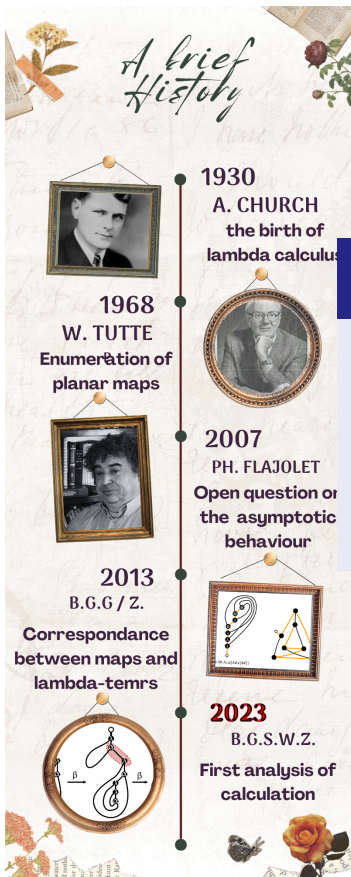
The pattern  $p_3$  is very hard to specify (perhaps impossible due to auto-correlation). But Alexandros Singh has proposed an elegant probabilistic method to get around this difficulty !

## Theorem (BGSWZ'2023)

*The mean number  $\mu_3$  of pattern  $p_3$  in closed linear terms of size  $n$  verifies asymptotically  $\mu_3 \geq \frac{n}{240}$*

In fact, we can have the equality, but this needs very sophisticated guess and proof approach based on differential algebra.

# Towards calculation on linear $\lambda$ -terms



## Theorem (BGSWZ'2023)

The mean length  $\ell_n$  of the  $\beta$ -reduction chain for a closed linear terms of size  $n$  verifies asymptotically

$$\ell_n \geq \frac{11n}{240}$$

Very closed to Noam's conjecture of  $n/21$   
 ( $1/21 - 11/240 = 0.0017\dots$ )

## Perspectives and (still) open questions

- Find the asymptotic of general lambda-terms (work in progress with H.K.)
- Almost nothing is known about parameters of general lambda-terms
- Prove or disprove Noam's conjecture on *beta*-reduction (In particular, we currently only have a lower bound)
- Have more informations on the size of the terms during the *beta*-reduction process...

MERCI

THANK YOU

謝謝