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# INFINITE-ARMS BANDIT: OPTIMALITY VIA CONFIDENCE BOUNDS

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*Abstract:* The literature on the infinite-arms bandit problem includes a regret lower bound of all allocation strategies for Bernoulli rewards with a uniform prior, and strategies based on success runs. Furthermore, a two-target algorithm has been proposed that achieves the regret lower bound, and optimality has been extended to Bernoulli rewards with general priors. We present a confidence-bound target (CBT) algorithm that achieves optimality for rewards that are bounded above. For each arm, we construct a confidence bound and compare it against those of other arms and a target value to determine whether the arm should be sampled further. The target value depends on the assumed priors of the arm means. In the absence of information on the prior, the target value is determined empirically. Numerical studies show that the CBT algorithm is versatile and outperforms its competitors.

*Key words and phrases:* MAB, optimal allocation, sequential analysis, UCB.

## 1. Introduction

Berry et al. (1997) initiated the development of the infinite-arms bandit problem. For Bernoulli rewards with a uniform prior, they showed a  $\sqrt{2n}$  regret lower bound for  $n$  rewards, and provided algorithms based on success runs that achieve no more than  $2\sqrt{n}$  regret. Bonald and Proutière (2013) provided a two-target stopping-time algorithm that can get arbitrarily close to the lower bound of Berry et al. (1997), and is also optimal on Bernoulli rewards with general priors. Wang, Audibert, and Munos (2008) considered bounded rewards, and showed that their confidence-bound algorithm has regret bounds that are  $\log n$  times the optimal regret. Vermorel and Mohri (2005) proposed a POKER algorithm for general reward distributions and priors.

The confidence-bound method has been arguably the most influential approach over the past 30 years for the fixed arm-size bandit problem. Lai and Robbins (1985) derived the smallest asymptotic regret that can be achieved by any algorithm. Lai (1987) showed that by constructing an upper confidence bound (UCB) for each arm, and playing the arm with the largest UCB, this smallest regret is achieved in exponential families. The UCB approach was subsequently extended to unknown time horizons and other parametric families in Agrawal (1995a), Auer, Cesa-Bianchi, and

Fischer (2002), Burnetas and Katehakis (1996), Cappé et al. (2013), and Kaufmann, Cappé, and Garivier (2012), and has been shown to perform well in practice, achieving optimality beyond that of exponential families. Chan (2020) modified the subsampling approach of Baransi, Maillard, and Mannor (2014) to show that optimality is achieved in exponential families, despite not applying parametric information in the selection of the arms. The method applies confidence bounds that are computed empirically from subsample information, which substitutes for the missing parametric information. A related problem is the study of the multi-armed bandit with irreversible constraints, initiated by Hu and Wei (1989).

The Bayesian approach has also enjoyed considerable success; see Berry and Fristedt (1985), Gittins (1989), and Thompson (1933) for early ground-work, and Korda, Kaufmann, and Munos (2013) for more recent advances.

We show here how the confidence-bound method can be applied on infinite arms. We call this new procedure the confidence-bound target (CBT). As in the UCB, in the CBT, a confidence bound is computed for each arm. The difference is that in the CBT, we specify an additional target value. Then, we compare the confidence bound of an arm against this target to decide whether to play the arm further, or to discard it and play a new arm.

We derive a regret lower bound that applies to all bandit algorithms, and proceed to show how the target in the CBT is chosen to achieve this lower bound. This optimal target depends only on the prior distribution of the arm means, and not on the reward distributions. That is, the reward distributions need not be specified for optimality to be achieved.

To handle the situation in which the prior is not available, we provide an empirical version of the CBT in which the target value is computed empirically. Numerical studies on Bernoulli rewards and on a URL data set show that the CBT and empirical CBT outperform their competitors.

In the related continuum-armed bandit problem, there is an uncountably infinite number of arms. Each arm is indexed by a known parameter  $\theta$  and has rewards with mean  $f(\theta)$ , where  $f$  is an unknown continuous function. For solutions to the problem of maximizing the expected sum of rewards, see Agrawal (1995b), Auer, Ortner, and Szepesvári (2007), Cope (2009), Kleinberg (2004), and Tyagi and Gärtner (2013).

The remainder of this paper proceeds as follows. In Section 2, we describe the infinite-arms bandit problem. In Section 3, we review the literature on this problem. In Section 4, we describe the CBT. In Section 5, we motivate why the chosen target of the CBT leads to the regret lower bound, and state the optimality of the CBT. In Section 6, we introduce an

empirical version of the CBT to tackle unknown priors, and explain why it works. In Section 7, we perform numerical studies. Section 8 concludes the paper.

## 2. Problem setup

Let  $Y_{k1}, Y_{k2}, \dots$  be independent and identically distributed (i.i.d.) rewards from arm  $k$ . In the classical multi-armed bandit problem, there are finitely many arms, and the objective is to sequentially select the arms so as to maximize the expected sum of rewards. Equivalently, we minimize the regret, which is the expected cumulative difference between the best arm mean and the mean of the arm played.

In the infinite-arms bandit problem that we consider here, there are infinitely many arms, and rewards are bounded above by a value that we shall assume for simplicity to be one. We assume, in addition, that it is possible for an arm to have reward mean arbitrarily close to one.

The regret of a bandit algorithm, after  $n$  trials, is defined to be

$$R_n = E\left(\sum_{k=1}^{\infty} \sum_{t=1}^{n_k} X_{kt}\right), \text{ where } X_{kt} = 1 - Y_{kt} (\geq 0) \quad (2.1)$$

is the loss associated with reward  $Y_{kt}$ , and  $n_k$  is the number of times arm  $k$  has been played (hence,  $n = \sum_{k=1}^{\infty} n_k$ ). The expectation in (2.1) is with respect to the following Bayesian framework.

Let  $g$  be a prior on  $(0, \infty)$ . For each  $\mu$  in which  $g(\mu) > 0$ , let  $F_\mu$  be a nonnegative distribution with mean  $\mu$ . The expectation in (2.1) is with respect to

$$\mu_k \stackrel{\text{i.i.d.}}{\sim} g \text{ for } k \geq 1 \text{ and } X_{kt} \stackrel{\text{i.i.d.}}{\sim} F_{\mu_k} \text{ for } t \geq 1. \quad (2.2)$$

The minimization of the regret under (2.2) for finite arms, known as the *stochastic bandit problem*, has been studied in Agrawal and Goyal (2012), Bubeck and Liu (2013), and Russo and Van Roy (2014).

In the infinite-arms bandit problem, a key decision to be made at each trial is whether to sample a new arm or to play a previously played arm. The Bayesian framework in (2.2) provides useful information on the new arms.

### 3. Preliminary background

Let  $a \wedge b$  denote  $\min(a, b)$ ,  $\lfloor \cdot \rfloor$  ( $\lceil \cdot \rceil$ ) denote the greatest (least) integer function, and  $a^+$  denote  $\max(0, a)$ . We say that  $a_n \sim b_n$  if  $\lim_{n \rightarrow \infty} (a_n/b_n) = 1$ ,  $a_n = o(b_n)$  if  $\lim_{n \rightarrow \infty} (a_n/b_n) = 0$ , and  $a_n = O(b_n)$  if  $\limsup_{n \rightarrow \infty} |a_n/b_n| < \infty$ .

Berry et al. (1997) showed that for Bernoulli rewards with  $g$  uniform on  $(0, 1)$ , a regret lower bound

$$\liminf_{n \rightarrow \infty} \frac{R_n}{\sqrt{n}} \geq \sqrt{2} \quad (3.1)$$

is unavoidable. They proposed the following bandit strategies:

1.  $f$ -failure strategy. We play the same arm until  $f$  failures are encountered. When this happens, we switch to a new arm. We do not go back to a previously played arm; that is, the strategy is *non-recalling*.
2.  $s$ -run strategy. We restrict ourselves to no more than  $s$  arms, following the one-failure strategy in each, until a success run of length  $s$  is observed in an arm. When this happens, we play the arm for the remaining trials. If no success run of length  $s$  is observed in all  $s$  arms, then the arm with the highest proportion of success is played for the remaining trials.
3. Non-recalling  $s$ -run strategy. We follow the one-failure strategy until an arm produces a success run of length  $s$ . When this happens, we play the arm for the remaining trials. If no arm produces a success run of length  $s$ , then the one-failure strategy is used in all  $n$  trials.
4.  $m$ -learning strategy. We follow the one-failure strategy for the first  $m$  trials, with the arm at trial  $m$  played until it yields a failure. Thereafter, we play, for the remaining trials, the arm with the highest proportion of successes.



Berry et al. (1997) showed that  $R_n \sim n/(\log n)$  for the  $f$ -failure strategy for any  $f \geq 1$ , whereas for the  $\sqrt{n}$ -run strategy, the  $\sqrt{n} \log n$ -learning strategy, and the non-recalling  $\sqrt{n}$ -run strategy,

$$\limsup_{n \rightarrow \infty} \frac{R_n}{\sqrt{n}} \leq 2.$$

Bonald and Proutière (2013) proposed a two-target algorithm with target values  $s_1 = \lfloor \sqrt[3]{\frac{n}{2}} \rfloor$  and  $s_f = \lfloor f\sqrt{\frac{n}{2}} \rfloor$ , where  $f \geq 2$  is user-defined. An arm is discarded if it has fewer than  $s_1$  successes when it encounters its first failure, or fewer than  $s_f$  successes when it encounters its  $f$ th failure. If both targets are met, then the arm is accepted and played for the remaining trials. Bonald and Proutière (2013) showed that for the uniform prior, the two-target algorithm satisfies, for  $n \geq \frac{f^2}{2}$ ,

$$R_n \leq f + \left(\frac{s_f+1}{f}\right) \left(\frac{s_f-f+2}{s_f-s_1-f+2}\right)^f \left(2 + \frac{1}{f} + \frac{2(f+1)}{s_1+1}\right),$$

from which they conclude that

$$\limsup_{n \rightarrow \infty} \frac{R_n}{\sqrt{n}} \leq \sqrt{2} + \frac{1}{f\sqrt{2}}.$$

Thus, for  $f$  and  $n$  large, the regret is close to the asymptotic lower bound  $\sqrt{2n}$ .

Bonald and Proutière (2013) extended their optimality on Bernoulli rewards to nonuniform priors. They showed that when  $g(\mu) \sim \alpha\mu^{\beta-1}$ , for

some  $\alpha > 0$  and  $\beta > 0$  as  $\mu \rightarrow 0$ , the regret lower bound of Berry et al. (1997) is extended to

$$\liminf_{n \rightarrow \infty} (n^{-\frac{\beta}{\beta+1}} R_n) \geq C_0, \text{ where } C_0 = \left(\frac{\beta(\beta+1)}{\alpha}\right)^{\frac{1}{\beta+1}}. \quad (3.2)$$

They also showed that their two-target algorithm with  $s_1 = \lfloor n^{\frac{1}{\beta+2}} C_0^{-\frac{\beta+1}{\beta+2}} \rfloor$  and  $s_f = \lfloor f n^{\frac{1}{\beta+1}} C_0^{-1} \rfloor$  satisfies

$$\limsup_{f \rightarrow \infty} [\limsup_{n \rightarrow \infty} (n^{-\frac{\beta}{\beta+1}} R_n)] \leq C_0.$$

Wang, Audibert, and Munos (2008) proposed a UCB-F algorithm for rewards taking values in  $[0, 1]$ , and showed that under suitable regularity conditions,  $R_n = O(n^{\frac{\beta}{\beta+1}} \log n)$ . In the UCB-F, an order  $n^{\frac{\beta}{\beta+1}}$  arms are chosen, and confidence bounds are computed on these arms to determine which arm to play. The UCB-F is different from the CBT in that it pre-selects the number of arms, and also does not have a mechanism to reject weak arms quickly. Carpentier and Valko (2015) also considered rewards taking values in  $[0, 1]$ , but their interest in maximizing the selection of a good arm differs from the aims here and in the papers above.

#### 4. Proposed methodology

We propose a new bandit algorithm called CBT, in which a confidence bound is constructed for each arm and compared against a target value.

Let  $S_{kt} = \sum_{u=1}^t X_{ku}$ ,  $\bar{X}_{kt} = t^{-1}S_{kt}$ , and  $\hat{\sigma}_{kt}^2 = t^{-1} \sum_{u=1}^t (X_{ku} - \bar{X}_{kt})^2$ . Let  $b_n$  and  $c_n$  be positive confidence coefficients satisfying

$$b_n \rightarrow \infty \text{ and } c_n \rightarrow \infty \text{ with } b_n + c_n = o(n^\delta), \text{ for all } \delta > 0. \quad (4.1)$$

In our numerical studies, we select  $b_n = c_n = \log \log n$ . We define the confidence bound of arm  $k$ , after it has been played  $t$  times, to be

$$L_{kt} = \max \left( \frac{\bar{X}_{kt}}{b_n}, \bar{X}_{kt} - \frac{c_n \hat{\sigma}_{kt}}{\sqrt{t}} \right). \quad (4.2)$$

Let  $\zeta > 0$  be a specified target value. In the CBT, the arms are played sequentially. Arm  $k$  is played until  $L_{kt}$  goes above  $\zeta$ , and it is discarded when that happens. We discuss in Section 5 how  $\zeta$  should be selected to achieve optimality. It suffices to mention here that the optimal  $\zeta$  decreases at a polynomial rate with respect to  $n$ .

#### Confidence bound target (CBT)

1. Play arm 1 at trial 1.
2. For  $m = 1, \dots, n - 1$ : Let  $k$  be the arm played at trial  $m$ , and let  $t$  be the number of times arm  $k$  has been played up to trial  $m$ .
  - (a) If  $L_{kt} \leq \zeta$ , then play arm  $k$  at trial  $m + 1$ .

(b) If  $L_{kt} > \zeta$ , then play arm  $k + 1$  at trial  $m + 1$ .

Let  $K$  be the number of arms played after  $n$  trials, and let  $n_k$  be the number of times arm  $k$  has been played after  $n$  trials. Hence,  $n = \sum_{k=1}^K n_k$ .

There are three types of arms that we need to take care of, and that explains the design of  $L_{kt}$ . The first type is arms with  $\mu_k$  (mean of loss  $X_{kt}$ ) significantly larger than  $\zeta$ . We would like to reject these arms quickly. The decision to reject arm  $k$  when  $\bar{X}_{kt}/b_n$  exceeds  $\zeta$  is key to achieving this.

The second type is arms with  $\mu_k$  larger than  $\zeta$ , but not by as much as those of the first type. We are unlikely to reject these arms quickly because it is difficult to determine whether  $\mu_k$  is smaller or larger than  $\zeta$  based on a small sample. Rejecting arm  $k$  when  $\bar{X}_{kt} - c_n \hat{\sigma}_{kt}/\sqrt{t}$  exceeds  $\zeta$  ensures that arm  $k$  is rejected only when it is statistically significant that  $\mu_k$  is larger than  $\zeta$ . Though there may be a large number of rewards from these arms, their contributions to the regret are small because they have small  $\mu_k$ , because  $\zeta$  is chosen small when  $n$  is large.

The third type of arms is those with  $\mu_k$  smaller than  $\zeta$ . For these arms, the best strategy (when  $\zeta$  is chosen correctly) is to play them for the remaining trials. Selecting  $b_n \rightarrow \infty$  and  $c_n \rightarrow \infty$  in (4.2) ensures that the probabilities of rejecting these arms are small.

For Bernoulli rewards, the first target  $s_1$  of the two-target algorithm is designed for the quick rejection of the first type of arms, and the second target  $s_f$  is designed to reject the second type. The difference between the two-target algorithm and the CBT is that whereas the former monitors an arm for rejection only when there are one and  $f$  failures, with  $f$  chosen large for optimality, the CBT checks for rejection each time a failure occurs. The frequent monitoring by the CBT is a positive feature that results in significantly better performance in the numerical experiments discussed in Section 7.

## 5. Optimality

We state the regret lower bound in Section 5.1, and show that the CBT achieves this bound in Section 5.2.

### 5.1 Regret lower bound

In Lemma 1, we motivate the choice of  $\zeta$ . Let  $P_\mu$  denote the probability and  $E_\mu$  denote the expectation with respect to  $X \stackrel{d}{\sim} F_\mu$ . Let  $P_g(\cdot) = \int_0^\infty P_\mu(\cdot)g(\mu)d\mu$  and  $E_g(\cdot) = \int_0^\infty E_\mu(\cdot)g(\mu)d\mu$ . Let  $\lambda = \int_0^\infty E_\mu(X|X > 0)g(\mu)d\mu [= E_g(X|X > 0)]$  be the mean of the first positive loss of a random arm. We assume that  $\lambda < \infty$ . The value  $\lambda$  is the unavoidable cost of

exploring a new arm. We consider  $E_\mu(X|X > 0)$  instead of  $\mu$  because it makes sense to reject an arm only after observing a positive loss. For Bernoulli rewards,  $\lambda = 1$ . Let  $p(\zeta) = P_g(\mu_1 \leq \zeta)$  and  $v(\zeta) = E_g(\zeta - \mu_1)^+$ .

Consider an idealized algorithm that plays arm  $k$  until a positive loss is observed, and  $\mu_k$  is revealed when that happens. If  $\mu_k > \zeta$ , then arm  $k$  is rejected, and a new arm is played next. If  $\mu_k \leq \zeta$ , then we stop exploring and play arm  $k$  for the remaining trials.

Let

$$r_n(\zeta) = \frac{\lambda}{p(\zeta)} + nE_g(\mu_1 | \mu_1 \leq \zeta). \quad (5.1)$$

Assuming that the exploration stage of the idealized algorithm uses  $o(n)$  trials and  $\zeta$  is small, its regret is asymptotically  $r_n(\zeta)$ . Let  $K$  be the total number of arms played. The first term in the expansion of  $r_n(\zeta)$  approximates  $E(\sum_{k=1}^{K-1} \sum_{t=1}^{n_k} X_{kt})$ , whereas the second term approximates  $E(\sum_{t=1}^{n_K} X_{Kt})$ .

**Lemma 1.** *Let  $\zeta_n$  be such that  $v(\zeta_n) = \frac{\lambda}{n}$ . We have*

$$\inf_{\zeta > 0} r_n(\zeta) = r_n(\zeta_n) = n\zeta_n.$$

PROOF. Because  $E_g(\zeta - \mu_1 | \mu_1 \leq \zeta) = \frac{v(\zeta)}{p(\zeta)}$ , it follows from (5.1) that

$$r_n(\zeta) = \frac{\lambda}{p(\zeta)} + n\zeta - \frac{nv(\zeta)}{p(\zeta)}. \quad (5.2)$$

It follows from  $\frac{d}{d\zeta}v(\zeta) = p(\zeta)$  and  $\frac{d}{d\zeta}p(\zeta) = g(\zeta)$  that

$$\frac{d}{d\zeta}r_n(\zeta) = \frac{g(\zeta)[nv(\zeta) - \lambda]}{p^2(\zeta)}.$$

Because  $v$  is continuous and strictly increasing when it is positive, the root to  $v(\zeta) = \frac{\lambda}{n}$  exists, and Lemma 1 follows from solving  $\frac{d}{d\zeta}r_n(\zeta) = 0$ .  $\square$

Consider:

(A1) there exist  $\alpha > 0$  and  $\beta > 0$  such that  $g(\mu) \sim \alpha\mu^{\beta-1}$  as  $\mu \rightarrow 0$ .

Under (A1),  $p(\zeta) = \int_0^\zeta g(\mu)d\mu \sim \frac{\alpha}{\beta}\zeta^\beta$  and  $v(\zeta) = \int_0^\zeta p(\mu)d\mu \sim \frac{\alpha}{\beta(\beta+1)}\zeta^{\beta+1}$  as  $\zeta \rightarrow 0$ ; hence,  $v(\zeta_n) \sim \frac{\lambda}{n}$  for

$$\zeta_n \sim Cn^{-\frac{1}{\beta+1}}, \text{ where } C = \left(\frac{\lambda\beta(\beta+1)}{\alpha}\right)^{\frac{1}{\beta+1}}. \quad (5.3)$$

In Lemma 2, we state the regret lower bound. We assume the following:

(A2) there exists  $a_1 > 0$  such that  $P_\mu(X > 0) \geq a_1 \min(\mu, 1)$ , for all  $\mu$ .

We need this assumption to avoid having bad arms that are played a large number of times, because their losses are mostly zeros, but can be very big when positive.

**Lemma 2.** *Under (A1) and (A2), all infinite-arms bandit algorithms have regret satisfying*

$$R_n \geq [1 + o(1)]n\zeta_n \sim Cn^{\frac{\beta}{\beta+1}} \text{ as } n \rightarrow \infty. \quad (5.4)$$

Lemma 2 is proved in the Supplementary Material.

**EXAMPLE 1.** Consider  $X \stackrel{d}{\sim} \text{Bernoulli}(\mu)$ . Condition (A2) holds with  $a_1 = 1$ . If  $g$  is uniform on  $(0,1)$ , then (A1) holds with  $\alpha = \beta = 1$ . Because

$\lambda = 1$ , by (5.3),  $\zeta_n \sim \sqrt{\frac{2}{n}}$ . Lemma 2 states that  $R_n \geq [1 + o(1)]\sqrt{2n}$ , agreeing with Theorem 3 of Berry et al. (1997).

Bonald and Proutière (2013) showed (5.4) in their Lemma 3 for Bernoulli rewards under (A1), and showed that their two-target algorithm gets close to the regret lower bound when  $f$  is large. We show in Theorem 1 that the lower bound in (5.4) is achieved by the CBT for rewards that need not be Bernoulli.

## 5.2 Optimality of the CBT

We state the optimality of the CBT in Theorem 1, after describing the conditions on discrete rewards under (B1) and continuous rewards under (B2) for which the theorem holds. Let  $M_\mu(\theta) = E_\mu e^{\theta X}$ .

(B1) The rewards are integer-valued. For  $0 < \delta \leq 1$ , there exists  $\theta_\delta > 0$  such that for  $\mu > 0$  and  $0 \leq \theta \leq \theta_\delta$ ,

$$M_\mu(\theta) \leq e^{(1+\delta)\theta\mu}, \quad (5.5)$$

$$M_\mu(-\theta) \leq e^{-(1-\delta)\theta\mu}. \quad (5.6)$$

In addition,

$$P_\mu(X > 0) \leq a_2\mu \text{ for some } a_2 > 0, \quad (5.7)$$

$$E_\mu X^4 = O(\mu) \text{ as } \mu \rightarrow 0. \quad (5.8)$$



(B2) The rewards are continuous random variables satisfying

$$\sup_{\mu>0} P_\mu(X \leq \gamma\mu) \rightarrow 0 \text{ as } \gamma \rightarrow 0. \quad (5.9)$$

Moreover, (5.8) holds and for  $0 < \delta \leq 1$ , there exists  $\tau_\delta > 0$  such that for  $0 < \theta\mu \leq \tau_\delta$ ,

$$M_\mu(\theta) \leq e^{(1+\delta)\theta\mu}, \quad (5.10)$$

$$M_\mu(-\theta) \leq e^{-(1-\delta)\theta\mu}. \quad (5.11)$$

In addition, for each  $t \geq 1$ , there exists  $\xi_t > 0$  such that

$$\sup_{\mu \leq \xi_t} P_\mu(\hat{\sigma}_t^2 \leq \gamma\mu^2) \rightarrow 0 \text{ as } \gamma \rightarrow 0, \quad (5.12)$$

where  $\hat{\sigma}_t^2 = t^{-1} \sum_{u=1}^t (X_u - \bar{X}_t)^2$  and  $\bar{X}_t = t^{-1} \sum_{u=1}^t X_u$  for i.i.d.  $X_u \stackrel{d}{\sim} F_\mu$ .

**Theorem 1.** *Assume (A1), (A2), and either (B1) or (B2). For the CBT with threshold  $\zeta_n$  satisfying (5.3) and  $b_n, c_n$  satisfying (4.1),*

$$R_n \sim n\zeta_n \text{ as } n \rightarrow \infty. \quad (5.13)$$

Theorem 1 states that the CBT is optimal because it attains the lower bound given in Lemma 2. In the examples below, we show that the regularity conditions (A2), (B1), and (B2) are reasonable and checkable. The proof of Theorem 1 and the verification details in Examples 3–5 are given in the Supplementary Material.

EXAMPLE 2. If  $X \stackrel{d}{\sim} \text{Bernoulli}(\mu)$  under  $P_\mu$ , then

$$M_\mu(\theta) = 1 - \mu + \mu e^\theta \leq \exp[\mu(e^\theta - 1)].$$

Hence, (5.5) and (5.6) hold with  $\theta_\delta > 0$  satisfying

$$e^{\theta_\delta} - 1 \leq \theta_\delta(1 + \delta) \text{ and } e^{-\theta_\delta} - 1 \leq -\theta_\delta(1 - \delta).$$

In addition, (5.7) holds with  $a_2 = 1$ , and (5.8) holds because  $E_\mu X^4 = \mu$ .

Condition (A2) holds with  $a_1 = 1$ .

EXAMPLE 3. Let  $F_\mu$  be a distribution with support on  $0, \dots, I$  for some positive integer  $I > 1$  and having mean  $\mu$ . Condition (A2) holds with  $a_1 = I^{-1}$ , and (B1) holds as well.

EXAMPLE 4. Let  $F_\mu$  be the Poisson distribution with mean  $\mu$ . Condition (A2) holds with  $a_1 = 1 - e^{-1}$ , and (B1) holds as well.

EXAMPLE 5. Let  $Z$  be a continuous nonnegative random variable with mean one, and with  $Ee^{\tau_0 Z} < \infty$ , for some  $\tau_0 > 0$ . Let  $F_\mu$  be the distribution of  $\mu Z$ . Condition (A2) holds with  $a_1 = 1$ , and (B2) holds as well.

## 6. Methodology for unknown priors

The optimal implementation of the CBT and, in particular, the computation of the optimal target  $\zeta_n$ , assumes knowledge of how  $g(\mu)$  behaves for

$\mu$  near zero. For  $g$  unknown, we rely on Theorem 1 to motivate an empirical implementation of the CBT.

What is striking about (5.13) is that it relates the optimal target  $\zeta_n$  to  $\frac{R_n}{n}$ ; moreover, this relation does not depend on either the prior  $g$  or the reward distributions. We suggest, therefore, in an empirical implementation of the CBT, to apply the targets

$$\zeta(m) := \frac{S'_m}{n}, \quad (6.1)$$

where  $S'_m$  is the sum of the losses  $X_{kt}$  over the first  $m$  trials.

In the beginning with  $m$  small,  $\zeta(m)$  underestimates the optimal target, but this encourages exploration, which is the right strategy at the beginning. As  $m$  increases,  $\zeta(m)$  gets closer to the optimal target, and the empirical CBT behaves like the CBT when deciding whether to play an arm further. A key difference between the CBT and the empirical CBT is that the latter decides from among all played arms which to play further, whereas the CBT plays the arms sequentially.

### Empirical CBT

Notation: When there are  $m$  total rewards, let  $n_k(m)$  denote the number of rewards from arm  $k$ , and let  $K_m$  denote the number of arms played.

For  $m = 0$ , play arm 1. Hence,  $K_1 = 1$ ,  $n_1(1) = 1$ , and  $n_k(1) = 0$  for  $k > 1$ .

For  $m = 1, \dots, n - 1$ :

1. If  $\min_{1 \leq k \leq K_m} L_{kn_k(m)} \leq \zeta(m)$ , then play the arm  $k$  minimizing  $L_{kn_k(m)}$  at trial  $m + 1$ .
2. If  $\min_{1 \leq k \leq K_m} L_{kn_k(m)} > \zeta(m)$ , then play a new arm  $K_m + 1$  at trial  $m + 1$ .

The empirical CBT, unlike the CBT, does not achieve the smallest regret. This is because when a good arm (i.e., an arm with  $\mu_k$  below the optimal target) appears early, we are not sure whether this is because of good fortune or because the prior is disposed toward arms with small  $\mu_k$ . Thus, we explore more arms before we are certain, and play the good arm for the remaining trials. Similarly, when no good arm appears after many trials, we conclude that the prior is disposed toward arms with large  $\mu_k$ , and play an arm with  $\mu_k$  above the optimal target for the remaining trials, even though it is advantageous to explore further.

Because analyzing the regret of the empirical CBT is complicated, we consider an idealized version of the empirical CBT in the Supplementary

Material, deriving its asymptotic regret, to give us a sense of the additional regret when applying the CBT empirically.

In the idealized version of the empirical CBT,  $\mu_k$  is revealed after the first positive loss of arm  $k$  is observed. The number of arms played is the smallest  $K$  satisfying

$$\min_{1 \leq k \leq K} \mu_k \leq \frac{K\lambda}{n},$$

and exploitation of the best arm begins after  $\mu_1, \dots, \mu_K$  have been revealed. The idealized empirical CBT is like the idealized algorithm described at the beginning of Section 5.1, but with a target  $\zeta = \frac{k\lambda}{n}$ , after  $k$  arms have been played. This is because  $\lambda$  is the mean of the first positive loss of each arm, so after  $k$  arms have been played, the sum of the losses has mean  $k\lambda$ . The idealized empirical CBT is a simplification of the empirical CBT that captures the additional regret of the empirical CBT over that of the CBT when applying a target that does not depend on the prior.

**Theorem 2.** *The idealized empirical CBT has regret*

$$R'_n \sim I_\beta n \zeta_n, \tag{6.2}$$

where  $I_\beta = \left(\frac{1}{\beta+1}\right)^{\frac{1}{\beta+1}} \left(2 - \frac{1}{(\beta+1)^2}\right) \Gamma\left(2 - \frac{\beta}{\beta+1}\right)$  and  $\Gamma(u) = \int_0^\infty x^{u-1} e^{-x} dx$ .

The constant  $I_\beta$  increases with  $\beta$ , with  $I_0 = 1$  and  $\lim_{\beta \rightarrow \infty} I_\beta = 2$ . The increase is quite slow, so that for reasonable values of  $\beta$ , it is closer to one

than two, for example,  $I_1 = 1.10$ ,  $I_3 = 1.24$ , and  $I_5 = 1.36$ . Equation (6.2) states that the empirical CBT should have regret not more than 36% over the baseline lower bound when  $\beta \leq 5$ . This agrees with the simulation outcomes presented in Section 7.

## 7. Numerical studies

We study arms with Bernoulli rewards in Example 6, and arms with unspecified reward distributions in Example 7. In our simulations, 10,000 data sets are generated for each entry in Tables 1–4, and standard errors are placed after the  $\pm$  sign. In both the CBT and the empirical CBT, we select  $b_n = c_n = \log \log n$ . Aziz (2019) performed numerical studies involving various infinite-arms bandit algorithms, including the CBT and empirical CBT, with the objective of finding the arm with the best mean. The study also applies an infinite-arms bandit to an online data set involving voting responses to 3795 proposed captions for a cartoon on a New Yorker website.

**EXAMPLE 6.** We consider Bernoulli rewards with uniform prior  $g(\mu) = 1$ , as well as the beta priors  $g(\mu) = 3\mu^2$  [i.e., Beta(3,1)],  $g(\mu) = \frac{15}{16}\mu^2(1-\mu)^{-\frac{1}{2}}$  [i.e., Beta(3,  $\frac{1}{2}$ )],  $g(\mu) = 5\mu^4$  [i.e., Beta(5,1)], and  $g(\mu) = \frac{315}{256}\mu^4(1-\mu)^{-\frac{1}{2}}$  [i.e., Beta(5,  $\frac{1}{2}$ )].

We see from Tables 1–3 that the two-target algorithm does better with

		Regret			
		$n = 100$	$n = 1000$	$n = 10,000$	$n = 100,000$
CBT	$\zeta = \sqrt{2/n}$	14.6±0.1	51.5±0.3	162±1	504±3
	empirical	15.6±0.1	54.0±0.3	172±1	531±3
Berry et al.	1-failure	21.8±0.1	152.0±0.6	1123±4	8955±28
	$\sqrt{n}$ -run	19.1±0.2	74.7±0.7	260±3	844±9
	$\sqrt{n}$ -run (non-recall)	15.4±0.1	57.7±0.4	193±1	618±4
	$n^{\frac{1}{2}} \log n$ -learning	18.7±0.1	84.4±0.6	311±3	1060±9
Two-target	$f = 3$	15.2±0.1	52.7±0.3	167±1	534±3
	$f = 6$	16.3±0.1	55.8±0.4	165±1	511±3
	$f = 9$	17.5±0.1	58.8±0.4	173±1	514±3
UCB-F	$K = \lfloor \sqrt{n/2} \rfloor$	39.2±0.1	206.4±0.4	1204±1	4432±5
Lower bound	$\sqrt{2n}$	14.1	44.7	141	447

Table 1: Regrets for Bernoulli rewards with uniform prior.

$f = 3$  at smaller  $n$ , and with  $f = 6$  or  $9$  at larger  $n$ . The CBT is the best performer uniformly over sample size and prior, and the empirical CBT is competitive against the two-target algorithm, with  $f$  fixed.

Even though the CBT outperforms the empirical CBT, its optimal target  $\zeta$  depends on the prior. On the other hand, when applying the empirical CBT, the same algorithm is used for all priors here and on the URL data set in Example 7 with an unspecified prior. Hence, though it seems that

		Regret			
		Beta(3, 1)	Beta(3, $\frac{1}{2}$ )	Beta(5, 1)	Beta(5, $\frac{1}{2}$ )
CBT	$\zeta = Cn^{-\frac{1}{\beta+1}}$	284.2±0.9	363.0±1.0	474.0±1.0	554.3±1.1
	empirical	299.6±0.9	382.2±1.1	509.6±1.0	592.7±1.0
$n^{\frac{1}{\beta+1}}$ -run	non-recall	346.3±1.3	445.7±1.7	546.5±1.4	658.8±1.6
Two-target	$f = 3$	310.7±1.1	390.8±1.3	510.3±1.3	592.1±1.3
	$f = 6$	301.2±1.2	385.9±1.4	520.9±1.5	619.5±1.6
	$f = 9$	311.0±1.3	400.1±1.6	545.3±1.6	649.6±1.7
UCB-F		649.5±0.3	779.2±0.3	774.0±0.3	867.6±0.2
Lower bound	$Cn^{\frac{\beta}{\beta+1}}$	251.5	336.4	426.3	538.5

Table 2: Regrets for Bernoulli rewards with beta priors at  $n = 1000$ .

the empirical CBT is numerically comparable to the two-target algorithm and inferior to the CBT, in applications where the prior is unknown or incorrectly specified, it can perform much better.

For the uniform prior, the best performing among the algorithms in Berry et al. (1997) is the non-recalling  $\sqrt{n}$ -run algorithm. For the UCB-F [cf. Wang, Audibert, and Munos (2008)], the selection of  $K = \lfloor (\frac{\beta}{\alpha})^{\frac{1}{\beta+1}} (\frac{n}{\beta+1})^{\frac{\beta}{\beta+1}} \rfloor$  ( $\sim \frac{1}{p(\zeta_n)}$ ) and the “exploration sequence”  $\mathcal{E}_m = \sqrt{\log m}$  works well.

EXAMPLE 7. We consider the URL data set studied in Vermorel and Mohri (2005), who propose a POKER algorithm for dealing with a large number of arms. We reproduce part of their Table 1 in our Table 4, to-



		Regret ( $\times 10$ )			
		Beta(3, 1)	Beta(3, $\frac{1}{2}$ )	Beta(5, 1)	Beta(5, $\frac{1}{2}$ )
CBT	$\zeta = Cn^{-\frac{1}{\beta+1}}$	866 $\pm$ 3	1127 $\pm$ 4	2122 $\pm$ 5	2569 $\pm$ 6
	empirical	1004 $\pm$ 3	1318 $\pm$ 4	2547 $\pm$ 5	3149 $\pm$ 6
$n^{\frac{1}{\beta+1}}$ -run	non-recall	1476 $\pm$ 7	1713 $\pm$ 8	3874 $\pm$ 13	4142 $\pm$ 13
Two-target	$f = 3$	1159 $\pm$ 5	1501 $\pm$ 6	2973 $\pm$ 9	3559 $\pm$ 11
	$f = 6$	990 $\pm$ 4	1308 $\pm$ 5	2527 $\pm$ 7	3060 $\pm$ 9
	$f = 9$	957 $\pm$ 4	1257 $\pm$ 5	2429 $\pm$ 7	2992 $\pm$ 9
UCB-F		3739 $\pm$ 3	4522 $\pm$ 4	6488 $\pm$ 4	7499 $\pm$ 5
Lower bound	$Cn^{\frac{\beta}{\beta+1}}$	795	1064	1979	2499

Table 3: Regrets ( $\times 10$ ) for Bernoulli rewards with beta priors at  $n=100,000$ .

		Regret	
Algorithm	$\epsilon$	$n = 130$	$n = 1300$
emp. CBT		212 $\pm$ 2	123.8 $\pm$ 0.6
POKER		203	132
$\epsilon$ -greedy	0.05	733	431
$\epsilon$ -first	0.15	725	411
$\epsilon$ -decreasing	1.0	738	411

Table 4: Average regret  $R_n/n$ .

gether with new simulations on the empirical CBT. The data set consists of the retrieval latency of 760 university home pages, in milliseconds, with

a sample size of more than 1300 for each home page. The numbers in the data set correspond to the nonnegative losses  $X_{kt}$ . The data set can be downloaded from “sourceforge.net/projects/bandit.”

In our simulations, the losses are randomly permuted within home page in each run. At  $n = 130$ , POKER performs better than the empirical CBT, whereas at  $n = 1300$ , the empirical CBT performs better. The other algorithms are uniformly worse than both the POKER and the empirical CBT.

The algorithm  $\epsilon$ -first refers to exploring the first  $\epsilon n$  losses, with a random selection of the arms to be played. This is followed by pure exploitation for the remaining  $(1 - \epsilon)n$  losses, on the “best” arm (with the smallest mean loss). The algorithm  $\epsilon$ -greedy refers to selecting, in each trial, a random arm with probability  $\epsilon$ , and the best arm with the remaining  $1 - \epsilon$  probability. The algorithm  $\epsilon$ -decreasing is like  $\epsilon$ -greedy, except that in the  $m$ th trial, we select a random arm with probability  $\min(1, \frac{\epsilon}{m})$ , and the best arm otherwise. Both  $\epsilon$ -greedy and  $\epsilon$ -decreasing are disadvantaged by not using information on the total number of trials. Vermorel and Mohri (2005) also ran simulations on more complicated strategies, such as LeastTaken, SoftMax, Exp3, GaussMatch, and IntEstim, with the average regret ranging from 276–747 at  $n = 130$  and 189–599 at  $n = 1300$ .

## 8. Conclusion

The CBT optimizes the regret in the infinite-arms bandit problem when it is possible for an arm to have reward mean arbitrarily close to the upper bound of the rewards. This optimality is over all bandit algorithms, and does not require knowledge of the reward distribution for a given arm mean. It depends, however, on the correct selection of a target value computed from an assumed prior.

The empirical CBT is like the CBT, with the key difference being that the former computes the target value empirically. Though not optimal, it performs well in numerical studies, and is more practical because it can be applied without assuming a prior.

We suggest here two extensions of the CBT and empirical CBT for future work. The first is to handle the situation in which the sample size is not known in advance. Bonald and Proutière (2013) have a version of the two-target algorithm that they believe to be optimal for Bernoulli rewards when the sample size is not known in advance.

The second extension is to incorporate covariate information in the computation of the confidence bounds, leading to recommended arms that are specific to subgroups of the population. Modern developments in the finite-arms bandit literature has centered on handling covariate information;

see, for example, Goldenshluger and Zeevi (2013), Perchet and Rigollet (2013), Slivkins (2014), Wang, Kulkarni, and Poor (2005), and Yang and Zhu (2002). When the number of arms is comparable to or larger than the sample size, an infinite-arms approach is more appropriate and will provide strategies that differ from those of a finite-arms framework.

### Supplementary Material

The proofs of Lemma 2 and Theorems 1 and 2, and the verifications of (A2), (B1), and (B2) in Examples 3–5 are provided in the online Supplementary Material.

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